

PARALLEL ARCHITECTURES

Course “Parallel Computing”



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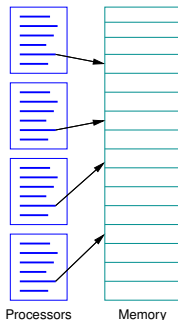
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Parallel Random Access Machine (PRAM)

- A simple abstract machine model.
 - Arbitrarily many processors execute same program on single shared memory.
 - Processors run synchronously in “lock-step”, possibly on different memory locations.
 - Cost of accessing memory is $O(1)$.
- Read/write conflicts have to be resolved.
 - EREW (exclusive read, exclusive write).
 - CREW (concurrent read, exclusive write).
 - CRCW (concurrent read, concurrent write) with multiple modes: common (only same value may be written), arbitrary (random value is written), priority (value of lowest numbered processor is written), . . .

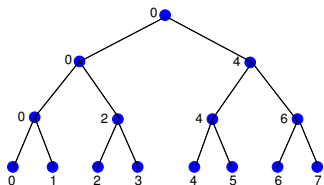


Traditional model for the analysis of parallel algorithms.

A PRAM Program

Multiply two matrices A and B of dimension $n = 2^m$.

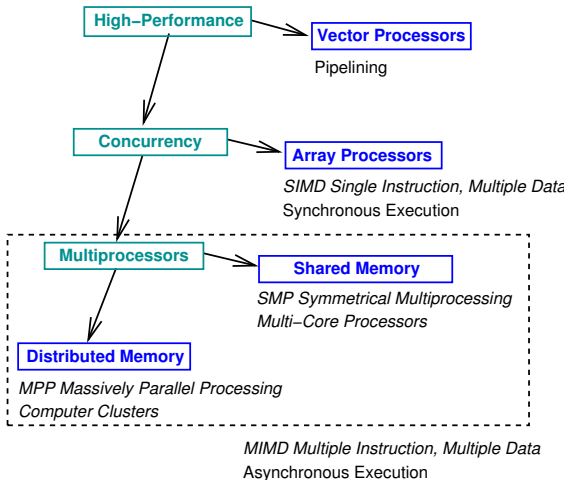
```
for i=0 to n-1 do in parallel
  for j=0 to n-1 do in parallel
    for k=0 to n-1 do in parallel
      P[i,j,k] = A[i,k]*B[k,j]
    for s=0 to m-1 do
      for t=0 to n-1 do in parallel
        if t%(2^(s+1)) = 0 then
          P[i,j,t] = P[i,j,t]+P[i,j,t+2^s]
      C[i,j] = P[i,j,0]
```



- First compute $P[i, j, k] = A[i, k] \times B[k, j]$ in time $O(1)$.
- Then compute $C[i, j] = \sum_k P[i, j, k]$ in time $O(\log n)$.
 - Computation of sum by $\log n$ stages of pairwise additions.

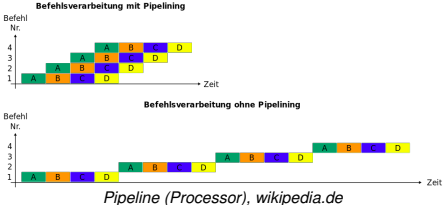
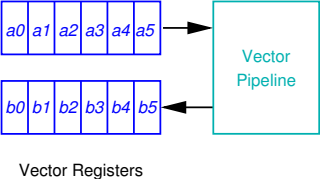
Matrix multiplication in time $O(\log n)$ with $O(n^3)$ processors.

High-Performance Architectures



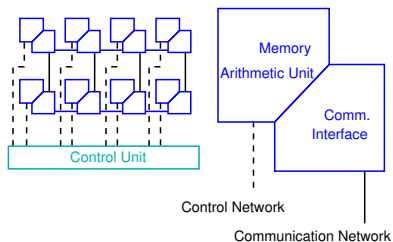
Flynn's Classification: SIMD versus MIMD.

Vector Processors



Vectors of data are processed by pipelines; speedup is limited by the (fixed) pipeline depth.

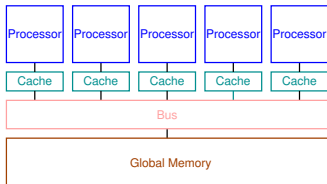
Array Processors



```
plural double matrix_multiply(A, B)
  plural double A,B;
{
  int i;
  plural double C = 0.0;
  ...
  for (i=0; i<nxproc; i++) {
    C += A*B;
    xnetW[1].A = A;
    xnetN[1].B = B;
  }
  return C;
}
```

Array of arithmetic units operates in lock-step.

Shared Memory Multi-Processors

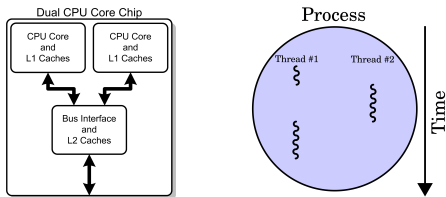


Alternative Term: SMP (Symmetric Multiprocessing)

- Multiple asynchronously operating processors.
- Single OS image schedules processes to processors.
- Single shared memory accessible via central bus.
 - Only one processor at a time can read/write memory.
 - Processors connected to bus via *coherent* caches.
 - *Snooping protocol*: whenever a cache sees another processor's write, it updates its local cache copies.

Scalable to 16 processors or so.

Multi-Core Processors

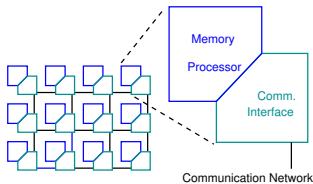


Multi-core processor, Thread, en.wikipedia.org

- Processors hold multiple processing units (“cores”).
 - Each core has a separate Level 1 cache.
 - Cores share a common Level 2 cache.
- Cores may execute multiple threads independently.
 - Threads: light-weight processes that can be independently scheduled for execution.
 - Processes: containers that hold multiple threads that have access to the same memory.

Today, actually every processor is by itself an SMP system.

Distributed Memory Multi-Processors

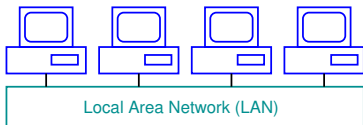


Alternative Term: MPP (Massively Parallel Processing)

- Many identical nodes that operate asynchronously.
 - Processor, local memory, communication interface.
- Each node runs its own OS image.
 - New processes are scheduled to the local processor.
- Nodes connected by high-bandwidth/low-latency network.
 - Different topologies (grid, tree, hypercube, ...).
 - Different network technologies (InfiniBand, OmniPath, ...)
 - Remote processes can communicate by *message passing*.

Scalable to thousands of processors.

Computer Clusters



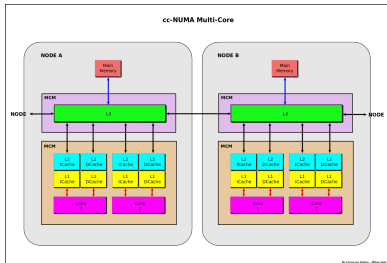
Also called “Beowulf” systems.

- A cheap alternative to MPP systems.
 - Each node is an independent of-the-shelf computer.
 - Connectivity provided by conventional Ethernet (or dedicated high-speed, e.g., InfiniBand) connections.
- A software stack implements MPP capabilities.
 - MPI, cluster managers, workload schedulers, ...
- Special programming and data processing software.
 - Apache Hadoop/MapReduce, Apache Spark, ...

Also MPPs are based on cluster-technology, so categories blur.

Virtual Shared Memory Multi-Processors

- ccNUMA: “cache coherent non-uniform memory access”.
 - All local memories combined to single address space.
 - NUMA: access to remote memory is more expensive.
 - Directory keeps track of which nodes hold cache copies of which lines of local memory.
 - If local memory line is updated, nodes with copies are informed.



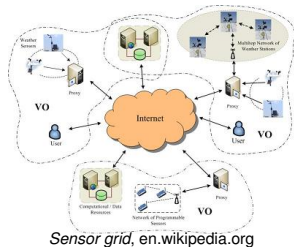
Cache memory, en.wikipedia.org

Implementation of SMP model on top of MPP hardware.

Computational Grids

Infrastructures composed of resources from multiple networks.

- Heterogenous combination of various kinds of resources.
 - Computing power, data storage, sensors, . . .
 - Located in different administrative domains.
- Connected by “Grid middleware”.
 - Globus Toolkit, gLite, Unicore, . . .



Software-based implementation of a (widely distributed) “virtual supercomputer”.

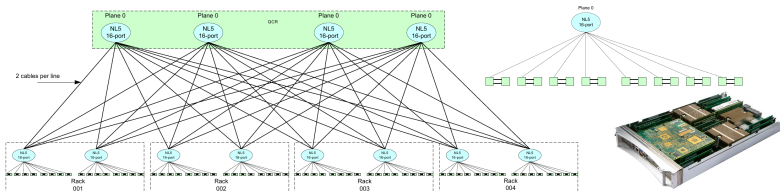
The JKU Supercomputer "Mach"

SGI UV-1000 an der Johannes Kepler Universität

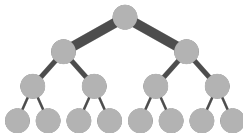
- Rechnerarchitektur
- Prozessortyp
- Prozessoranzahl
- Speicher
- Betriebssystem
- Prozessorleistung

- Memory-Bandbreite
- Bisection-Bandbreite

SGI UV-1000 shared Memory/cc-numa Architektur
Intel E78837 (Westmere - EX)
X86-64, 2.66GHz / 8-Cores / 24MB Cache
256 (2048 Cores)
16 TB shared Memory
Linux – Suse SLES 11 mit SGI Performance Suite
gesamt Peak = 21,3 TFlops
Spec_2006_INT Rate = ~39.000
Spec_2006_FP Rate = ~29.000
Stream = 5,8 Tbyte/s
Linpack 100 = ~2,2 Gflop/s
Linpack NxN = 18,5 Tflop/s
7,5 TB / s
480 GB / s

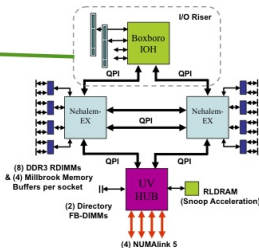


The Interconnection Topology “Fat Tree”



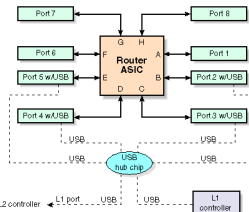
Fat tree, en.wikipedia.org

- Data links nearer to the top have higher bandwidth.
- Approximately same bisection bandwidth at each level.
- Implemented by proprietary NUMalink[®] technology.



Ports 1, 6, 7, and 8 can connect to compute or MPX modules

Ports 2, 3, 4, and 5 can connect to compute, MPX, or PCI expansion modules



The Data Access Hierarchy

Data Hierarchy Layer	Latency	Normalized Access Time
L1 Cache	1.4 ns	1×
L3 Cache	23 ns	16×
Local Memory	75 ns	53×
Remote Memory	1 μ s	700×
Disk	2 ms	3.6·10 ⁶ ×

Processors	Cores	Router Hops
2	16	0
32	256	1
256	2048	3

Considering the placement of processes and data is important for achieving high performance on a NUMA system.

Our Course Machine

Actually, we are going to use Mach's "little brother" Zusie.

- Configurations

- 1 blade: $2 \times 8 = 16$ cores, 128 GB RAM, 24 MB cache.
- Mach: $4 \times 32 = 128$ blades, 2048 cores, 16 TB RAM.
- Zusie: 32 blades, 512 cores, 2 TB RAM.

- Login

```
ssh -X -l ... zusie.jku.at
```

- Only from JKU network.
- Another possibility for use from other networks.

Be considerate: this machine is shared by many users.

Architecture Information

```
zusie> topology
Serial number: UV-00000044
Partition number: 0
    32 Blades
    1024 CPUs
2002.99 Gb Memory Total
    64.00 Gb Max Memory on any blade
    0.00 Gb Partition Base Address
4031.98 Gb Partition Last Address
    4 I/O Risers
    1 InfiniBand Controller
    4 Network Controllers
    2 SCSI Controllers
    8 USB Controllers
    1 VGA GPU
```

```
zusie> cat /proc/cpuinfo
processor      : 0
vendor_id    : GenuineIntel
cpu family   : 6
model        : 46
model name   : Intel(R) Xeon(R) CPU
stepping     : 6
cpu MHz      : 2267.156
cache size   : 24576 KB
....
processor     : 1023
vendor_id    : GenuineIntel
...
```

Hyper-Threading: 2 virtual cores per physical core,
thus 1024 virtual cores in total.

User and Process Information

```
zusie> who
```

```
root      console      2016-07-11 09:06
ws98      pts/0        2016-09-30 12:29
hans      pts/2        2017-02-21 17:04
hans      pts/3        2017-02-22 08:24 (:2.0)
hans      pts/4        2017-01-13 14:31 (lilli.edvz.uni-linz.ac.at)
k313270   pts/5        2017-02-22 08:38 (amir.risc.uni-linz.ac.at)
k313270   pts/6        2017-02-22 08:42 (amir.risc.uni-linz.ac.at)
...
```

```
zusie> ps -fu k313270
```

UID	PID	PPID	C	STIME	TTY	TIME	CMD
k313270	82369	1	0	Feb13	?	00:00:00	ssh -N -L 9999:localhost:37
k313270	447710	447708	0	08:38	?	00:00:00	sshd: k313270@pts/5
k313270	447711	447710	0	08:38	pts/5	00:00:00	-bash
k313270	449736	447800	99	08:50	pts/6	00:00:01	ps -fu k313270

Displays all users and all processes running on your behalf.

Thread Information

```
zusie> top -H -u k313270
```

```
top - 08:52:17 up 226 days, 52 min, 21 users,  load average: 214.24, 227.65, 24
```

```
Tasks: 17374 total, 218 running, 17155 sleeping,  1 stopped,  0 zombie
```

```
Cpu(s): 91.5%us,  5.7%sy,  2.3%ni,  0.3%id,  0.1%wa,  0.0%hi,  0.1%si,  0.0%st
```

```
Mem:   2051061M total, 1288263M used,  762798M free,          0M buffers
```

```
Swap:  131071M total,    568M used,  130503M free,  1119275M cached
```

PID	USER	PR	NI	VIRT	RES	SHR	S	%CPU	%MEM	TIME+	COMMAND
449747	k313270	20	0	62600	15m	1880	R	77	0.0	0:02.51	top
82369	k313270	20	0	60672	3432	2532	S	0	0.0	0:00.11	ssh
447710	k313270	20	0	106m	2540	1476	S	0	0.0	0:00.00	sshd
447711	k313270	20	0	55824	5684	2760	S	0	0.0	0:00.11	bash
447799	k313270	20	0	106m	2556	1476	S	0	0.0	0:00.16	sshd
447800	k313270	20	0	55824	5756	2800	S	0	0.0	0:00.28	bash

Displays all threads that are running on your behalf.

CPU Sets

```
zusie> cpuset -i /Upper256sh
```

```
zusie> my_cpuset.csh
```

```
cpuset: /Upper256sh
```

```
allowed resource Ids:
```

```
Cpus_allowed_list: 256-511
```

```
Mems_allowed_list: 32-63
```

```
zusie> jkutoptop -s /Upper256sh
```

```
...
```

```
zusie> man jkutoptop
```

```
...
```

```
JKUtop is a rewrite of top geared towards speed, especially on big SMP systems. ... It also has no functionality to display individual threads. Thanks to these measures it achieves a noticeable speedup on ridiculously large SMP systems compared to top from procs.
```

```
...
```

Creates new shell that confines processes to certain nodes.