VL Formal Modeling (Summer semester 2024)

Symbolic Summation and the modeling of sequences

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General picture:



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Part 1: Symbolic summation (a short introduction)

Part 2: Modeling of sequences with a term algebra (user interface)

Part 3: Modeling of sequences in difference rings (computer algebra)

Part 4: Construction of appropriate difference rings (advanced CA)

Part 5: Applications

Part 1: Symbolic Summation (a short introduction)

We start with the following telescoping problem:

Given an expression f(k) that evaluates to a sequence. **Find** an expression g(k) such that the telescoping equation holds:

$$f(k) = g(k+1) - g(k)$$
(1)

Suppose we find such an expression g(k). Then we proceed as follows. Summing (1) over k from a to b (and assuming that no poles arise during the evaluation) yields

k = 1

$$\sum_{k=a}^{b} f(k) = g(b+1) - g(a).$$
(2)

Note: we could always choose

$$g(k) = \sum_{i=a}^{n-1} f(i) \tag{3}$$

which would turn (2) to the trivial identity $\sum_{k=a}^{b} f(k) = \sum_{k=a}^{b} f(k)$. Thus we should refine our problem from above:

Find an expression g(k) with (1) where g(k) is simpler than the trivial solution (3).

Indefinite summation of polynomials

We start with one of the most simplest cases: the summand is a polynomial, i.e., $f(x) \in \mathbb{K}[x]$.

The following questions arise:

- 1. What is the domain of expressions in which we search g(k)?
- 2. How can we calculate a solution g(k) in this solution domain?

As it turns out, the first question can be answered nicely: a solution g(x) exists always in $\mathbb{K}[x]$. For the second question, we will consider two different tactics that are often used in summation packages.

Tactic 1: the classical approach

Note that for indefinite integration of polynomials one can utilize the following well known property: for any $m \in \mathbb{N}$ we have

$$D_x x^m = m \, x^{m-1}$$

which implies

$$\int_{a}^{b} x^{m} dx = \frac{x^{m+1}}{m+1} \Big|_{a}^{b} = \frac{b^{m+1} - a^{m+1}}{m+1}.$$

Thus by linearity we can integrate any polynomial by

$$\int_{a}^{b} \sum_{m=0}^{d} c_{m} x^{m} dx = \sum_{m=0}^{d} c_{m} \int_{a}^{b} x^{m} dx = \sum_{m=0}^{d} \frac{c_{m}(b^{m+1} - a^{m+1})}{m+1}$$

For indefinite summation of polynomials we can follow precisely the same classical strategy.

Definition. For any sequence (expression) g(k) we define

$$\Delta g(k) := g(k+1) - g(k).$$

Lemma

For $m \in \mathbb{N}$ we have

$$\Delta x^{\underline{m}} = m \, x^{\underline{m-1}} \, .$$

Proof.

We have

$$\begin{aligned} \Delta x^{\underline{m}} &= (x+1)^{\underline{m}} - x^{\underline{m}} \\ &= (x+1)x(x-1)\dots(x-m+2) - x(x-1)\dots(x-m+1) \\ &= ((x+1) - (x-m+1))x(x-1)\dots(x-m+2) \\ &= m \, x^{\underline{m-1}}. \end{aligned}$$

As a consequence we get

$$\Delta \frac{x^{\underline{m+1}}}{\underline{m+1}} = x^{\underline{m}}, \quad m \in \mathbb{N}$$

and summing this equation over k from a to b yields

$$\sum_{x=a}^{b} x^{\underline{m}} = \frac{(b+1)^{\underline{m+1}} - a^{\underline{m+1}}}{m+1}$$

Note that this is nothing else than the continuous version for integration. In particular, for given d

$$f(x) = \sum_{m=0}^{u} c_m \, x^{\underline{m}} \in \mathbb{K}[x]$$

with $d \in \mathbb{N}$ it follows that

$$g(x) = \sum_{m=0}^{a} \frac{c_m x^{\underline{m+1}}}{m+1}$$

is a telescoping solution. Furthermore,

$$\sum_{x=a}^{b} f(x) = \sum_{m=0}^{d} c_m \sum_{k=a}^{b} k^{\underline{m}} = \sum_{m=0}^{d} \frac{c_m((b+1)^{\underline{m+1}} - a^{\underline{m+1}})}{m+1}.$$

Part 1: Symbolic Summation (a short introduction)

The only problem is that in many cases one does not have a polynomial given in the representation of falling factorials but in the standard form

$$\sum_{m=0}^{d} \bar{c}_m \, x^m \in \mathbb{K}[x].$$

Luckily one can rewrite a polynomial written in the basis

$$1, x, x^2, \ldots, x^d$$

to the representation written in the basis

$$x^{\underline{0}} = 1, x^{\underline{1}} = x, x^{\underline{2}} = x(x-1), \dots, x^{\underline{d}} = x(x-1)\dots(x-d+1)$$

by using the formula

$$x^m = \sum_{k=0}^m S(m,k) x^{\underline{k}}$$

where $S(\boldsymbol{n},\boldsymbol{k})$ denotes the Stirling numbers of second kind. They can be computed by

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{i} \binom{k}{i} (k-i)^{n};$$

Example. Consider the polynomial

$$f(x) = x^4.$$

Using the formulas from above, we get

$$f(x) = x^{4} = \sum_{k=0}^{4} S(4,k)x^{\underline{k}} = 0x^{\underline{0}} + 1x^{\underline{1}} + 7x^{\underline{2}} + 6x^{\underline{3}} + 1x^{\underline{4}}.$$

Thus we get

$$g(x) = \frac{1}{2}x^2 + \frac{7}{3}x^3 + \frac{3}{2}x^4 + \frac{1}{5}x^5$$

= $\frac{1}{30}(x-1)x(2x-1)(3x^2-3x-1).$

such that

$$g(x+1) - g(x) = f(x)$$

holds. In particular we get

$$\sum_{k=1}^{n} k^4 = \sum_{k=1}^{n} f(k) = g(n+1) - g(1) = \frac{1}{30}n(n+1)(2n+1)(3n^2 + 3n - 1).$$

Tactic 2: linear algebra.

We use the following property: for $f(x) \in \mathbb{K}[x]$ there is a $g(x) \in \mathbb{K}[x]$ with (1) where

$$\deg(g) \le \deg(f) + 1.$$

Thus setting $d:= \deg(f)+1$ for given $f \in \mathbb{K}[x]$ the desired solution has the form

$$g(x) = \sum_{m=0}^{d} g_m \, x^m$$

and we can determine the unknowns $g_0, \ldots, g_d \in \mathbb{K}$ by linear algebra as follows.

Example. Take $f(x) = x^4 \in \mathbb{Q}[x]$. With $d = \deg(f) + 1 = 5$ the ansatz $g(x) = g_0 + g_1 x + g_2 x^2 + g_3 x^3 + g_4 x^4 + g_5 x^5$

for the unknowns $g_0, g_1, g_2, g_3, g_4, g_5 \in \mathbb{Q}$ is in place. This gives

$$\begin{aligned} x^4 = &\Delta g(x) = g(x+1) - g(x) = 0x^5 \\ &+ 5g_5 x^4 \\ &+ (4g_4 + 10g_5)x^3 \\ &+ (3g_3 + 6g_4 + 10g_5)x^2 \\ &+ (2g_2 + 3g_3 + 4g_4 + 5g_5)x \\ &+ (g_1 + g_2 + g_3 + g_4 + g_5)x^0. \end{aligned}$$

By coefficient comparison this yields the linear system

$$\begin{array}{ll} [x^4] & 1 = 5g_5 \\ [x^3] & 0 = 4g_4 + 10g_5 \\ [x^2] & 0 = 3g_3 + 6g_4 + 10g_5 \\ [x^1] & 0 = 2g_2 + 3g_3 + 4g_4 + 5g_5 \\ [x^0] & 0 = g_1 + g_2 + g_3 + g_4 + g_5 \end{array}$$

which is already in triangular form.

Thus we can read off the solution

$$g_5 = \frac{1}{5}, \quad g_4 = -\frac{1}{2}, \quad g_3 = \frac{1}{3}, \quad g_2 = 0, \quad g_1 = -\frac{1}{30}, \quad g_0 = c$$

with $c \in \mathbb{Q}$. In particular, we can choose c = 0 and obtain

$$g(x) = \frac{x^5}{5} - \frac{x^4}{2} + \frac{x^3}{3} - \frac{x}{30} = \frac{1}{30}(x-1)x(2x-1)\left(3x^2 - 3x - 1\right).$$

To this end, we continue as in the previous example and get the desired result.

More general summation objects for indefinite and definite summation

Clearly, the first tactic is very elegant, but it works only for the special case of polynomial summation. For the second tactic one has to work more (i.e., has to solve in addition a linear system), but it turns out to be more general. More precisely, one can carry over these ideas to a rather general setting that works not only for the polynomial ring $\mathbb{Q}[x]$ but in more general rings called $R\Pi\Sigma$ -difference rings that have been implemented within the summation package Sigma. In the following all technical details are omitted and we proceed with a concrete example.

Example. We want to sum

$$\sum_{k=0}^{n} H_k$$

In order to accomplish this task, we take

$$f(k) = H_k$$

and search for

$$g(k) \in \mathbb{Q}(k)[H_k]$$

with

$$f(k) = g(k+1) - g(k).$$

Here we can use a similar tactic as used in the case of polynomial summation. Namely, summation theory tells us that any such solution g(k) has the property

$$\deg(g) \le \deg(f) + 1 = 1 + 1 = 2.$$

As a consequence we can make the ansatz

$$g(k) = g_0(k)H_k^0 + g_1(k)H_k^1 + g_2(k)H_k^2$$

with $g_0(k), g_1(k), g_2(k) \in \mathbb{Q}(k)$.

Using recursive algorithms and linear system solving (details are skipped here) we find

$$g_0(k) = -k$$
$$g_1(k) = k$$
$$g_2(k) = 0,$$

i.e.,

$$g(k) = -k + kH_k + 0H_k^2.$$

Hence summing the telescoping equation over k from 0 to n gives

$$\sum_{k=0}^{n} H_k = g(n+1) - g(0) = (n+1)H_{n+1} - (n+1) = -n + (1+n)H_n.$$

The above machinery can be carried out within the summation package Sigma. After loading it into Mathematica

ln[1] = << Sigma.m

Sigma - A summation package by Carsten Schneider (C) RISC-JKU

one can insert the above sum

$$\begin{split} & \text{In}[2]:= \mathbf{mySum} = \mathbf{SigmaSum}[\mathbf{SigmaHNumber}[k], \{k, 0, 1\}] \\ & \text{Out}[2]= \sum_{k=0}^{n} \mathbb{H}_{k} \end{split}$$

and can apply the command

 $ln[3] := \mathbf{SigmaReduce}[\mathbf{mySum}]$

 $\texttt{Out[3]}= -n + (1+n)\texttt{H}_n$

In general one can insert, e.g., a sum of the form

$$\sum_{k=l}^{n} f(k)$$

with $l \in \mathbb{N}$ where f(k) itself is given in terms of indefinite nested sums defined over hypergeometric products.

Definition

Let \mathbb{K} be a field. A product $\prod_{j=l}^{k} f(j)$, $l \in \mathbb{N}$, is called **hypergeometric in** k over \mathbb{K} if $f(x) \in \mathbb{K}(x)$ is a rational function where the numerator and denominator of f(j) are nonzero for all $j \in \mathbb{Z}$ with $j \ge l$. An expression in terms indefinite of nested sums over hypergeometric products in k over \mathbb{K} is composed recursively by the three operations $(+, -, \cdot)$ with

- elements from the rational function field $\mathbb{K}(k)$,
- hypergeometric products in k over \mathbb{K} ,
- ▶ and sums of the form $\sum_{j=l}^{k} f(j)$ with $l \in \mathbb{N}$ where f(j) is an expression in terms of indefinite nested sums over hypergeometric products in j over \mathbb{K} ; here it is assumed that the evaluation of f(j) for all $j \geq l$ does not introduce any poles.

ln[4] := mySum =

$$\begin{split} \mathbf{SigmaSum}[\mathbf{SigmaPower}[-1,k]\mathbf{SigmaBinomial}[n,k]\mathbf{SigmaHNumber}[k],\{k,a,b\}\\ \mathsf{Out}[4]=& \sum_{k=a}^{b}(-1)^{k}\binom{n}{k}\mathsf{H}_{k} \end{split}$$

$$\begin{array}{l} \mbox{In[5]:= } \mathbf{SigmaReduce[mySum]} \\ \mbox{Out[5]:= } & \big(\frac{(a-n)(-1+a-n)}{an^2} + \frac{(-1+a-n) \mathbb{H}_a}{n} \big) (-1)^{1+a} \binom{n}{-1+a} + \big(\frac{-b+n}{n^2} + \frac{(-b+n) \mathbb{H}_b}{n} \big) (-1)^b \binom{n}{b} \end{array}$$

$$\begin{split} &\ln[6] = \mathbf{mySum} = \mathbf{SigmaSum}[\mathbf{SigmaBinomial}[n,k], \{k,0,r\}]^2, \{r,0,b\}] \\ &\text{Out}[6] = \sum_{r=0}^{b} \big(\sum_{k=0}^{r} \binom{n}{k}\big)^2 \end{split}$$

$$\begin{aligned} & \text{In[7]:= SigmaReduce[mySum]} \\ & \text{Out[7]:= } (-b+n)\binom{n}{b}\sum_{i_1=0}^{b}\binom{n}{i_1} + \frac{1}{2}(2+2b-n)\big(\sum_{i_1=0}^{b}\binom{n}{i_1}\big)^2 - \frac{1}{2}n\sum_{i_1=0}^{b}\binom{n}{i_1}^2 \end{aligned}$$

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▶ For any element $f = \frac{p}{q} \in \mathbb{G}$ with $p, q \in \mathbb{K}[x]$ where $q \neq 0$ and p, q being coprime we define

$$\operatorname{ev}(f,k) = \begin{cases} 0 & \text{if } q(k) = 0\\ \frac{p(k)}{q(k)} & \text{if } q(k) \neq 0. \end{cases}$$

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▶ We define L(f) to be the minimal value $\delta \in \mathbb{N}$ such that $q(k) \neq 0$ holds for all $k \geq \delta$; further,

$$Z(f) = \max(L(\frac{1}{p}), L(\frac{1}{q})) \quad \text{if } f \neq 0.$$

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We define L(f) to be the minimal value δ ∈ N such that q(k) ≠ 0 holds for all k ≥ δ; further,

$$Z(f) = \max(L(\frac{1}{p}), L(\frac{1}{q})) \quad \text{if } f \neq 0.$$

Example: For

$$f = \frac{p}{q} = \frac{x-4}{(x-3)(x-1)}$$

we get

$$(\operatorname{ev}(f,n))_{n\geq 0} = (-\frac{4}{3}, \underline{0}, 2, \underline{0}, 0, \frac{1}{8}, \dots) \in \mathbb{Q}^{\mathbb{N}}$$

For $n \ge L(f) = 4$ no poles arise; for $n \ge Z(f) = \max(L(\frac{1}{p}), L(\frac{1}{q})) = \max(4, 5) = 5$ no zeroes arise.

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We define

 $\mathcal{R} = \{r \in \mathbb{K} \setminus \{1\} \mid r \text{ is a root of unity}\}$

with the function $\operatorname{ord} : \mathcal{R} \to \mathbb{Z}_{\geq 1}$ where

$$\operatorname{ord}(r) = \min\{n \in \mathbb{Z}_{\geq 1} \mid r^n = 1\}.$$

Let $\otimes,\,\oplus,\,\odot,$ Sum, Prod and RPow be operations with the signatures

 $\mathbf{Prod}^*(\mathbb{G}) =$ the smallest set that contains 1 with the following properties:

1. If $r \in \mathcal{R}$ then $\mathsf{RPow}(r) \in \mathsf{Prod}^*(\mathbb{G})$.

Let $\otimes,\,\oplus,\,\odot,$ Sum, Prod and RPow be operations with the signatures

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Prod^{*}(\mathbb{G})= the smallest set that contains 1 with the following properties:

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- 3. If $p,q \in \mathsf{Prod}^*(\mathbb{G})$ then $p \odot q \in \mathsf{Prod}^*(\mathbb{G})$.

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 $\mathbf{Prod}^*(\mathbb{G})$ = the smallest set that contains 1 with the following properties:

If r ∈ R then RPow(r) ∈ Prod*(G).
 If f ∈ G* and l ∈ N with l ≥ Z(f) then Prod(l, f) ∈ Prod*(G).
 If p, q ∈ Prod*(G) then p ⊙ q ∈ Prod*(G).
 If p ∈ Prod*(G) and z ∈ Z \ {0} then p^Oz ∈ Prod*(G).

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Furthermore, we define

 $\Pi(\mathbb{G}) = \{ \mathsf{RPow}(r) \mid r \in \mathcal{R} \} \cup \{ \mathsf{Prod}(l, f) \mid f \in \mathbb{G}, l \in \mathbb{N} \}.$

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Example: In $\mathbb{G} = \mathbb{Q}(x)$ we get

$$P = (\underbrace{\mathsf{Prod}(1,x)}_{\in \Pi(\mathbb{G})} \overset{\textcircled{0}}{(-2)}) \odot \underbrace{\mathsf{RPow}(-1)}_{\Pi(\mathbb{G})} \in \mathsf{Prod}^*(\mathbb{G}).$$

$\textbf{SumProd}(\mathbb{G}) = \text{the smallest set containing } \mathbb{G} \cup \mathsf{Prod}^*(\mathbb{G})$ with:

1. For all $f,g \in \mathsf{SumProd}(\mathbb{G})$ we have $f \oplus g \in \mathsf{SumProd}(\mathbb{G})$.

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- 2. For all $f, g \in \mathsf{SumProd}(\mathbb{G})$ we have $f \odot g \in \mathsf{SumProd}(\mathbb{G})$.
- 3. For all $f \in \text{SumProd}(\mathbb{G})$ and $k \in \mathbb{Z}_{\geq 1}$ we have $f^{\otimes}k \in \text{SumProd}(\mathbb{G})$.

 $\textbf{SumProd}(\mathbb{G}) = \text{the smallest set containing } \mathbb{G} \cup \mathsf{Prod}^*(\mathbb{G})$ with:

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$\mathbb{G} \longrightarrow \mathsf{SumProd}(\mathbb{G}) \text{ (nested sums over hypergeometric products)}$

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Furthermore, the **set of nested sums over hypergeometric products** is given by

$$\Sigma(\mathbb{G}) = \{\mathsf{Sum}(l, f) \mid l \in \mathbb{N} \text{ and } f \in \mathsf{SumProd}(\mathbb{G})\}$$

and the set of nested sums and hypergeometric products is given by $\Sigma\Pi(\mathbb{G}) = \Sigma(\mathbb{G}) \cup \Pi(\mathbb{G}).$

$\mathbb{G} \longrightarrow \mathsf{SumProd}(\mathbb{G})$ (nested sums over hypergeometric products)

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Example

With $\mathbb{G} = \mathbb{K}(x)$ we get, e.g., the following expressions:

$$\begin{split} E_1 &= \mathsf{Sum}(1,\mathsf{Prod}(1,x)) \in \Sigma(\mathbb{G}) \subset \mathsf{SumProd}(\mathbb{G}), \\ E_2 &= \mathsf{Sum}(1,\frac{1}{x+1} \odot \mathsf{Sum}(1,\frac{1}{x^3}) \odot \mathsf{Sum}(1,\frac{1}{x})) \in \Sigma(\mathbb{G}) \subset \mathsf{SumProd}(\mathbb{G}), \\ E_3 &= (E_1 \oplus E_2) \odot E_1 \in \mathsf{SumProd}(\mathbb{G}). \end{split}$$

 $\operatorname{ev}: \mathbb{G} \times \mathbb{N} \to \mathbb{K} \longrightarrow \operatorname{ev}: \mathsf{SumProd}(\mathbb{G}) \times \mathbb{N} \to \mathbb{K}$

 $\mathrm{ev}:\mathbb{G}\times\mathbb{N}\to\mathbb{K}\qquad\longrightarrow\qquad\mathrm{ev}:\mathsf{SumProd}(\mathbb{G})\times\mathbb{N}\to\mathbb{K}$

1. For $f,g \in \mathsf{SumProd}(\mathbb{G})$, $k \in \mathbb{Z} \setminus \{0\}$ $(k > 0 \text{ if } f \notin \mathsf{Prod}^*(\mathbb{G}))$ we set

$$ev(f^{\otimes}k, n) := ev(f, n)^k,$$

$$ev(f \oplus g, n) := ev(f, n) + ev(g, n),$$

$$ev(f \odot g, n) := ev(f, n) ev(g, n);$$

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2. for $r \in \mathcal{R}$ and $Sum(l, f), Prod(\lambda, g) \in SumProd(\mathbb{G})$ we define

$$\begin{split} &\operatorname{ev}(\operatorname{\mathsf{RPow}}(r),n):=\prod_{i=1}^n r=r^n,\\ &\operatorname{ev}(\operatorname{\mathsf{Sum}}(l,f),n):=\sum_{i=l}^n\operatorname{ev}(f,i),\\ &\operatorname{ev}(\operatorname{\mathsf{Prod}}(\lambda,g),n):=\prod_{i=\lambda}^n\operatorname{ev}(g,i)=\prod_{i=\lambda}^ng(i). \end{split}$$

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Note: $\Pi(\mathbb{G})$ defines all hypergeometric products (which evaluate to sequences with non-zero entries).

f can be considered as a simple program and ev(f, n) with $n \in \mathbb{N}$ executes it (like an interpreter/compiler) yielding the nth entry of the represented sequence.

Definition

For $F \in \mathsf{SumProd}(\mathbb{G})$ and $n \in \mathbb{N}$ we write F(n) := ev(F, n).

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Example

For $E_i \in \mathsf{SumProd}(\mathbb{K}(x))$ with i = 1, 2, 3 we get

$$E_1(n) = \operatorname{ev}(E_1, n) = \operatorname{ev}(\operatorname{Sum}(1, \operatorname{Prod}(1, x)), n) = \sum_{k=1}^n \prod_{i=1}^k i = \sum_{k=1}^n k!,$$

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General picture:



An expression $A \in \mathsf{SumProd}(\mathbb{G})$ is in **reduced representation** if

$$A = (f_1 \odot P_1) \oplus (f_2 \odot P_2) \oplus \dots \oplus (f_r \odot P_r)$$
(4)

with $f_i \in \mathbb{G}^*$ and

$$P_i = (a_{i,1} \otimes z_{i,1}) \odot (a_{i,2} \otimes z_{i,2}) \odot \cdots \odot (a_{i,n_i} \otimes z_{i,n_i})$$

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for $1 \leq i \leq r$ with one of the three choices

▶ $a_{i,j} = \operatorname{Sum}(l_{i,j}, f_{i,j})$ with $l_{i,j} \in \mathbb{N}$, $f_{i,j} \in \operatorname{SumProd}(\mathbb{G})$ and $z_{i,j} \in \mathbb{Z}_{\geq 1}$, ▶ $a_{i,j} = \operatorname{Prod}(l_{i,j}, f_{i,j})$ with $l_{i,j} \in \mathbb{N}$, $f_{i,j} \in \operatorname{Prod}^*(\mathbb{G})$ and $z_{i,j} \in \mathbb{Z} \setminus \{0\}$, ▶ $a_{i,j} = \operatorname{RPow}(f_{i,j})$ with $f_{i,j} \in \mathcal{R}$ and $1 \leq z_{i,j} < \operatorname{ord}(r_{i,j})$

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and such that the following two properties hold:

- 1. for each $1 \leq i \leq r$ and $1 \leq j < j' < n_i$ we have $a_{i,j} \neq a_{i,j'}$;
- 2. for each $1 \leq i < i' \leq r$ with $n_i = n_j$ there does not exist a $\sigma \in S_{n_i}$ with $P_{i'} = (a_{i,\sigma(1)} \otimes z_{i,\sigma(1)}) \odot (a_{i,\sigma(2)} \otimes z_{i,\sigma(2)}) \odot \cdots \odot (a_{i,\sigma(n_i)} \otimes z_{i,\sigma(n_i)}).$

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 $H\in\mathsf{SumProd}(\mathbb{G})$ is in sum-product reduced representation if

- it is in reduced representation;
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 $E_3 = (E_1 \oplus E_2) \odot E_1$ is not in reduced representation

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Lemma

For any $A \in SumProd(\mathbb{G})$, there is a $B \in SumProd(\mathbb{G})$ in sum-product reduced representation and $\lambda \in \mathbb{N}$ such that

$$A(n) = B(n) \quad \forall n \ge \lambda.$$

SumProd (W, \mathbb{G}) =the set of elements from SumProd (\mathbb{G}) which are in reduced representation and the arising sums/products are taken from W.

$$\begin{split} \mathbf{SumProd}(W,\mathbb{G}) = & \text{the set of elements from SumProd}(\mathbb{G}) \text{ which} \\ & \text{are in reduced representation and the arising} \\ & \text{sums/products are taken from } W. \end{split}$$

▶ W is called **shift-closed over** G if for any $A \in \text{SumProd}(W, G)$, $s \in \mathbb{Z}$ there are $B \in \text{SumProd}(W, G)$ and $\delta \in \mathbb{N}$ such that

 $A(n+s) = B(n) \quad \forall n \ge \delta.$

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$$\begin{array}{ll} W \text{ is shift-stable} & \stackrel{\Rightarrow}{\not\leftarrow} & W \text{ is shift-closed} \\ & \not\leftarrow \end{array}$$

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$$A(n+s) = B(n) \quad \forall n \ge \delta.$$

- ▶ W is called **shift-stable over** G if for any product or sum in W the multiplicand or summand is built by sums and products from W.
- ▶ W is called **canonical reduced over** \mathbb{G} if for any $A, B \in \mathsf{SumProd}(W, \mathbb{G})$ with

$$A(n) = B(n) \quad \forall n \ge \delta$$

for some $\delta \in \mathbb{N}$ the following holds: A and B are the same up to permutations of the operands in \oplus and \odot .

- $W \subseteq \Sigma \Pi(\mathbb{G})$ is called σ -reduced over \mathbb{G} if
 - $1. \ {\rm the \ elements \ in \ } W$ are in sum-product reduced form,
 - 2. W is shift-stable (and thus shift-closed) and
 - 3. W is canonical reduced.

In particular, $A \in \mathsf{SumProd}(W, \mathbb{G})$ is called σ -reduced (w.r.t. W) if W is σ -reduced over \mathbb{G} .

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Problem SigmaReduce: Compute a σ -reduced representation

Given: $A_1, \ldots, A_u \in \mathsf{SumProd}(\mathbb{G})$ with $\mathbb{G} = \mathbb{K}(x)$. Find: a σ -reduced set $W = \{T_1, \ldots, T_e\} \subset \Sigma \Pi(\mathbb{G})$, $B_1 \ldots, B_u \in \mathsf{SumProd}(W, \mathbb{G})$ and $\delta_1, \ldots, \delta_u \in \mathbb{N}$ such that for all $1 \leq i \leq r$ we get

 $A_i(n) = B_i(n) \quad n \ge \delta_i.$

• Canonical representation in term algebras





 $\forall n \geq \delta \ \operatorname{ev}(A_1, n) = \operatorname{ev}(B_1, n)$

• Canonical representation in term algebras



• Canonical representation in term algebras



General picture:



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Part 1: Symbolic summation (a short introduction)

Part 2: Modeling of sequences with a term algebra (user interface)

Part 3: Modeling of sequences in difference rings (computer algebra)

Part 4: Construction of appropriate difference rings (advanced CA)

Part 5: Applications

Part 3: Modeling of sequences in difference rings (computer algebra)

Represent $H={\rm Sum}(1,\frac{1}{x})\in {\rm SumProd}(\mathbb{G})$ with $H(n)=H_n=\sum_{k=1}^n\frac{1}{k}.$

Part 3: Modeling of sequences in difference rings (computer algebra)

Represent $H = \text{Sum}(1, \frac{1}{x}) \in \text{SumProd}(\mathbb{G})$ with $H(n) = H_n = \sum_{k=1}^n \frac{1}{k}.$ 1. a formal ring $\mathbb{A} = \underbrace{\mathbb{Q}(x)}_{\substack{\text{rat. fu. field} \\ \text{polynomial ring}}} [s]$ Part 3: Modeling of sequences in difference rings (computer algebra)

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- 1. a formal ring $\mathbb{A} = \mathbb{Q}(x)[s]$
- 2. an evaluation function

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 $\mathrm{ev}: \quad \mathbb{Q}(x)[s] \times \mathbb{N} \quad \rightarrow \quad \mathbb{Q}$

 $\mathbf{ev}(\mathbf{s},\mathbf{n}) = \mathbf{H}_{\mathbf{n}}$

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$$ev: \quad \mathbb{Q}(x)[s] \times \mathbb{N} \quad \to \quad \mathbb{Q} \\ \left(\sum_{i=0}^{d} f_i s^i, n\right) \quad \mapsto \quad \sum_{i=0}^{d} ev'(f_i, n) H_n^i$$

 $ev(s, n) = H_n$

Definition: (\mathbb{A}, ev) is called an eval-ring

- 1. a formal ring $\mathbb{A} = \mathbb{Q}(x)[s]$
- 2. an evaluation function $\mathrm{ev}:\mathbb{A}\times\mathbb{N}\to\mathbb{Q}$

Consider the map

$$\begin{array}{rccc} \tau : & \mathbb{A} & \to & \mathbb{Q}^{\mathbb{N}} \\ & f & \mapsto & \langle \mathrm{ev}(f,n) \rangle_{n \geq 0} \end{array}$$

It is almost a ring homomorphism :

$$\tau(x)\tau(\frac{1}{x}) \qquad = \quad \langle 0,1,2,3,\dots\rangle\langle 0,1,\frac{1}{2},\frac{1}{3},\dots\rangle$$

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$$||$$
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$$\star$$
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Consider the map

$$\begin{array}{rccc} \tau : & \mathbb{A} & \to & \mathbb{Q}^{\mathbb{N}}/\sim \\ & f & \mapsto & \langle \mathrm{ev}(f,n) \rangle_{n \geq 0} \end{array}$$

 $(a_n) \sim (b_n)$ iff $a_n = b_n$ from a certain point on

It is a ring homomorphism :

$$\begin{aligned} \tau(x)\tau(\frac{1}{x}) &= \langle 0, 1, 2, 3, \dots \rangle \langle 0, 1, \frac{1}{2}, \frac{1}{3}, \dots \rangle \\ & & || \\ \langle 0, 1, 1, 1, \dots \rangle \\ & & || \\ \tau(x\frac{1}{x}) &= \tau(1) &= \langle 1, 1, 1, 1, \dots \rangle \end{aligned}$$

- 1. a formal ring $\mathbb{A} = \mathbb{Q}(x)[s]$
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Consider the map

$$\begin{array}{rcccc} \tau : & \mathbb{A} & \to & \mathbb{Q}^{\mathbb{N}}/\sim & & (a_n) \sim \\ & f & \mapsto & \langle \operatorname{ev}(f,n) \rangle_{n \geq 0} & & \text{from a} \end{array}$$

 $(a_n) \sim (b_n)$ iff $a_n = b_n$ from a certain point on

It is an injective ring homomorphism (ring embedding):

- Represent $H={\rm Sum}(1,\frac{1}{x})\in {\rm SumProd}(\mathbb{G})$ with $H(n)=H_n=\sum_{k=1}^n\frac{1}{k}.$
 - 1. a formal ring $\mathbb{A}=\mathbb{Q}(x)[s]$
 - 2. an evaluation function $\mathrm{ev}:\mathbb{A}\times\mathbb{N}\to\mathbb{Q}$
 - 3. a ring automorphism

$$\begin{array}{rccc} \sigma': & \mathbb{Q}(x) & \to & \mathbb{Q}(x) \\ & r(x) & \mapsto & r(x+1) \end{array}$$

- 1. a formal ring $\mathbb{A}=\mathbb{Q}(x)[s]$
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$$\sigma': \quad \mathbb{Q}(x) \quad \to \quad \mathbb{Q}(x)$$
$$r(x) \quad \mapsto \quad r(x+1)$$
$$\sigma: \quad \mathbb{Q}(x)[s] \quad \to \quad \mathbb{Q}(x)[s] \qquad \qquad s \mapsto s + \frac{1}{x+1}$$

$$H_{n+1} = H_n + \frac{1}{n+1}$$

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$$\sigma: \quad \mathbb{Q}(x)[s] \quad \to \quad \mathbb{Q}(x)[s] \qquad \qquad s \mapsto s + \frac{1}{x+1}$$

$$\sum_{i=0}^{d} f_i s^i \quad \mapsto \quad \sum_{i=0}^{d} \sigma'(f_i) \left(s + \frac{1}{x+1}\right)^i \qquad H_{n+1} = H_n + \frac{1}{n+1}$$

Definition: (\mathbb{A}, σ) with a ring \mathbb{A} and automorphism σ is called a difference ring; the set of constants is

$$\operatorname{const}_{\sigma} \mathbb{A} = \{ c \in \mathbb{A} \mid \sigma(c) = c \}$$

- 1. a formal ring $\mathbb{A} = \mathbb{Q}(x)[s]$
- 2. an evaluation function $\mathrm{ev}:\mathbb{A}\times\mathbb{N}\to\mathbb{Q}$
- 3. a ring automorphism $\sigma : \mathbb{A} \to \mathbb{A}$

$$ev(\sigma(s), n) = ev(s + \frac{1}{x+1}, n) = H_n + \frac{1}{n+1} = ev(s, n+1)$$

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$$\tau(\sigma(s)) = \langle 1, 1 + \frac{1}{2}, 1 + \frac{1}{2} + \frac{1}{3}, \dots \rangle = S(\langle 0, 1, 1 + \frac{1}{2}, \dots \rangle) = S(\tau(s))$$

shift operator

- 1. a formal ring $\mathbb{A}=\mathbb{Q}(x)[s]$
- 2. an evaluation function $\mathrm{ev}:\mathbb{A}\times\mathbb{N}\to\mathbb{Q}$
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$$\tau \text{ is an injective difference ring homomorphism:}$$

$$\begin{split} \mathbb{K}(x)[s] & \xrightarrow{\sigma} & \mathbb{K}(x)[s] \\ & \downarrow^{\tau} & = & \downarrow^{\tau} \\ & \mathbb{K}^{\mathbb{N}}/\sim \xrightarrow{S} & \mathbb{K}^{\mathbb{N}}/\sim \end{split}$$

- 1. a formal ring $\mathbb{A}=\mathbb{Q}(x)[s]$
- 2. an evaluation function $\mathrm{ev}:\mathbb{A}\times\mathbb{N}\to\mathbb{Q}$
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$$\tau \text{ is an injective difference ring homomorphism:}$$

$$(\mathbb{K}(x)[s],\sigma) \stackrel{\tau}{\simeq} \boxed{(\underbrace{\tau(\mathbb{Q}(x))}_{\mathsf{rat. seq.}}[\langle H_n \rangle_{n \ge 0}], S)}_{\mathsf{rat. seq.}} \le (\mathbb{K}^{\mathbb{N}}/\sim, S)$$

Summary: we rephrase $H \in \text{SumProd}(\mathbb{G})$ as element h in a formal difference ring. More precisely, we will design

- ▶ a ring \mathbb{A} with $\mathbb{A} \supseteq \mathbb{G} \supseteq \mathbb{K}$ in which H can be represented by $h \in \mathbb{A}$;
- ▶ an evaluation function $ev : \mathbb{A} \times \mathbb{N} \to \mathbb{K}$ such that H(n) = ev(h, n)holds for sufficiently large $n \in \mathbb{N}$;
- ▶ a ring automorphism $\sigma : \mathbb{A} \to \mathbb{A}$ which models H(n+1) with $\sigma(h)$.

a ring

 $\mathbb{A} := \mathbb{K}(x)$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\sigma(x) = x + 1$$

a ring

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Sk!=(k+1)k!

a ring

 $\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}]$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\label{eq:scalar} \begin{split} \sigma(x) &= x+1\\ \mathsf{Sk!}{=}(\mathsf{k}{+}1)\mathsf{k!} &\leftrightarrow & \sigma(p_1) = (x+1)p_1 \end{split}$$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}]$$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\sigma(x) = x + 1$$

hypergeometric $\leftrightarrow \sigma(p_1) = c$ products

$$a_1 p_1 \qquad a_1 \in \mathbb{K}(x)^*$$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}]$$

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$$\sigma(x) = x + 1$$

$$\sigma(p_1) = a_1 p_1 \qquad a_1 \in \mathbb{K}(x)^* \\ \sigma(p_2) = a_2 p_2 \qquad a_2 \in \mathbb{K}(x)^*$$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}]$$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\sigma(x) = x + 1$$

hypergeometric + products

$$\begin{array}{ll} \stackrel{\rightarrow}{\rightarrow} & \sigma(p_1) = a_1 \, p_1 & a_1 \in \mathbb{K}(x)^* \\ & \sigma(p_2) = a_2 p_2 & a_2 \in \mathbb{K}(x)^* \\ & \vdots & \\ & \sigma(p_e) = a_e p_e & a_e \in \mathbb{K}(x)^* \end{array}$$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z]$$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\begin{aligned} \sigma(x) &= x + 1 \\ \text{hypergeometric} &\leftrightarrow \sigma(p_1) = a_1 p_1 & a_1 \in \mathbb{K}(x)^* \\ \text{products} & \sigma(p_2) = a_2 p_2 & a_2 \in \mathbb{K}(x)^* \\ &\vdots \\ &\sigma(p_e) = a_e p_e & a_e \in \mathbb{K}(x)^* \\ (-1)^k &\leftrightarrow \sigma(\mathbf{z}) = -\mathbf{z} & \mathbf{z}^2 = \mathbf{1} \end{aligned}$$

a ring

 γ

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z]$$

with an automorphism where $\sigma(c) = c$ for all $c \in \mathbb{K}$ and where

$$\begin{split} \sigma(x) &= x + 1 \\ \text{hypergeometric} &\leftrightarrow & \sigma(p_1) = a_1 \, p_1 & a_1 \in \mathbb{K}(x)^* \\ \text{products} & & \sigma(p_2) = a_2 p_2 & a_2 \in \mathbb{K}(x)^* \\ &\vdots & \\ & & \sigma(p_e) = a_e p_e & a_e \in \mathbb{K}(x)^* \\ \gamma \text{ is a primitive } \lambda \text{th} & \gamma^{\mathbf{k}} &\leftrightarrow & \sigma(\mathbf{z}) = \gamma \, \mathbf{z} & \mathbf{z}^{\lambda} = \mathbf{1} \end{split}$$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z][s_1]$$

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$$\sigma(x) = x + 1$$
hypergeometric $\leftrightarrow \sigma(p_1) = a_1 p_1$

$$a_1 \in \mathbb{K}(x)^*$$

$$\sigma(p_2) = a_2 p_2$$

$$a_2 \in \mathbb{K}(x)^*$$

$$\vdots$$

$$\sigma(p_e) = a_e p_e$$

$$a_e \in \mathbb{K}(x)^*$$

$$\sigma(\mathbf{z}) = \gamma \mathbf{z}$$

$$\mathbf{z}^{\lambda} = \mathbf{1}$$

$$\mathbf{z}^{\lambda} = \mathbf{1}$$

 $H_{k+1} = H_k + \frac{1}{k+1} \quad \leftrightarrow \quad \sigma(s_1) = s_1 + \frac{1}{x+1}$

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z][s_1]$$

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a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2]$$

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a ring

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- with an automorphism where $\sigma(c) = c$ for all $c \in \mathbb{K}$ and where

 $\sigma(x) = x + 1$ $\leftrightarrow \quad \sigma(p_1) = a_1 p_1 \qquad a_1 \in \mathbb{K}(x)^* \\ \sigma(p_2) = a_2 p_2 \qquad a_2 \in \mathbb{K}(x)^*$ hypergeometric products $\sigma(p_e) = a_e p_e \qquad a_e \in \mathbb{K}(x)^*$ γ is a primitive λ th $\gamma^{\mathbf{k}}$ \leftrightarrow $\sigma(\mathbf{z}) = \gamma \, \mathbf{z}$ $\mathbf{z}^{\lambda} = \mathbf{1}$ root of unity (nested) sum $\leftrightarrow \sigma(s_1) = s_1 + f_1 \quad f_1 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z]$ $\sigma(s_2) = s_2 + f_2 \quad f_2 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1]$ $\sigma(s_3) = s_3 + f_3$ $f_3 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2]$

Definition (Evaluation function)

Take (\mathbb{A}, σ) with a subfield \mathbb{K} of \mathbb{A} with $\sigma|_{\mathbb{K}} = \mathrm{id}$.

1. ev : $\mathbb{A} \times \mathbb{N} \to \mathbb{K}$ is called **evaluation function** for (\mathbb{A}, σ) if for all $f, g \in \mathbb{A}, c \in \mathbb{K}$ and $l \in \mathbb{Z}$ there exists a $\lambda \in \mathbb{N}$ with

$$\forall n \ge \lambda : \operatorname{ev}(c, n) = c,$$
(5)

$$\forall n \ge \lambda : \operatorname{ev}(f+g,n) = \operatorname{ev}(f,n) + \operatorname{ev}(g,n),$$
 (6)

$$\forall n \ge \lambda : \operatorname{ev}(f g, n) = \operatorname{ev}(f, n) \operatorname{ev}(g, n),$$
(7)

$$\forall n \ge \lambda : \operatorname{ev}(\sigma^{l}(f), n) = \operatorname{ev}(f, n+l).$$
(8)

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(6)

$$\forall n \ge \lambda : \operatorname{ev}(f g, n) = \operatorname{ev}(f, n) \ \operatorname{ev}(g, n), \tag{7}$$

$$\forall n \ge \lambda : \operatorname{ev}(\sigma^{l}(f), n) = \operatorname{ev}(f, n+l).$$
(8)

2. $L : \mathbb{A} \to \mathbb{N}$ is called *o*-function if for any $f, g \in \mathbb{A}$ with $\lambda = \max(L(f), L(g))$ the properties (6) and (7) hold and for any $f \in \mathbb{A}$ and $l \in \mathbb{Z}$ with $\lambda = L(f) + \max(0, -l)$ property (8) holds.

Connection between SumProd(\mathbb{G}) and hypergeometric *APS*-extension

• Observation 1: Given $\{T_1, \ldots, T_e\} \subseteq \Sigma \Pi(\mathbb{G})$, one can construct a hypergeometric *APS*-extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with ev and *L* such that there are $a_1, \ldots, a_e \in \mathbb{E}$ and $\delta_1, \ldots, \delta_e$ with $\operatorname{ev}(a_i, n) = T_i(n)$.

Connection between $\mathsf{SumProd}(\mathbb{G})$ and hypergeometric $\mathit{APS}\text{-extension}$

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• Observation 2:

 (\mathbb{E}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ a hypergeometric APS-extension of (\mathbb{G}, σ) ev : $\mathbb{E} \times \mathbb{N} \to \mathbb{K}$, $L : \mathbb{E} \to \mathbb{N}$

$$\forall n \ge L(t_i) : \\ \operatorname{ev}(t_i, n) = T_i(n) \in \Sigma \Pi(\mathbb{G})$$

 $W = \{T_1, \ldots, T_e\} \subseteq \Sigma \Pi(\mathbb{G})$ is sum-product reduced and shift stable: sums/products in T_i are from $\{T_1, \ldots, T_{i-1}\}$.

In particular, if $f \in \mathbb{E} \setminus \{0\}$, then we can take the "unique" $0 \neq F \in \mathsf{SumProd}(\{T_1, \ldots, T_e\}, \mathbb{G})$ with $F(n) = \operatorname{ev}(f, n)$ for all $n \geq L(f)$.

Connection between $\mathsf{SumProd}(\mathbb{G})$ and hypergeometric APS -extension

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$$\forall n \ge L(t_i) : \\ \operatorname{ev}(t_i, n) = T_i(n) \in \Sigma \Pi(\mathbb{G})$$

 $W = \{T_1, \ldots, T_e\} \subseteq \Sigma \Pi(\mathbb{G})$ is sum-product reduced and shift stable: sums/products in T_i are from $\{T_1, \ldots, T_{i-1}\}$.

In particular, if $f \in \mathbb{E} \setminus \{0\}$, then we can take the "unique" $0 \neq F \in \mathsf{SumProd}(\{T_1, \ldots, T_e\}, \mathbb{G})$ with $F(n) = \operatorname{ev}(f, n)$ for all $n \geq L(f)$. Definition

For $f \in \mathbb{E}$ we also write $\exp(f) = F$ for this particular F.

Connection between $\mathsf{SumProd}(\mathbb{G})$ and hypergeometric $\mathit{APS}\text{-extension}$

• Observation 1: Given $\{T_1, \ldots, T_e\} \subseteq \Sigma\Pi(\mathbb{G})$, one can construct a hypergeometric *APS*-extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with ev and *L* such that there are $a_1, \ldots, a_e \in \mathbb{E}$ and $\delta_1, \ldots, \delta_e$ with $\operatorname{ev}(a_i, n) = T_i(n)$.

• Observation 2:

 (\mathbb{E}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ a hypergeometric APS-extension of (\mathbb{G}, σ) ev : $\mathbb{E} \times \mathbb{N} \to \mathbb{K}$, $L : \mathbb{E} \to \mathbb{N}$

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Example

For
$$f = x + \frac{x+1}{x}s^4 \in \mathbb{Q}(x)[s]$$
 we obtain
 $\exp(f) = F = x \oplus (\frac{x+1}{x} \odot (\operatorname{Sum}(1, \frac{1}{x})^{\textcircled{0}}4) \in \operatorname{Sum}(\mathbb{Q}(x)))$
with $F(n) = \operatorname{ev}(f, n)$ for all $n \ge 1$.

Connection between $\mathsf{SumProd}(\mathbb{G})$ and hypergeometric APS -extension

• Observation 1: Given $\{T_1, \ldots, T_e\} \subseteq \Sigma\Pi(\mathbb{G})$, one can construct a hypergeometric *APS*-extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with ev and *L* such that there are $a_1, \ldots, a_e \in \mathbb{E}$ and $\delta_1, \ldots, \delta_e$ with $\operatorname{ev}(a_i, n) = T_i(n)$.

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 (\mathbb{E}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ a hypergeometric APS-extension of (\mathbb{G}, σ) ev : $\mathbb{E} \times \mathbb{N} \to \mathbb{K}$, $L : \mathbb{E} \to \mathbb{N}$

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 $W = \{T_1, \ldots, T_e\} \subseteq \Sigma \Pi(\mathbb{G})$ is sum-product reduced and shift stable: sums/products in T_i are from $\{T_1, \ldots, T_{i-1}\}$.

Difference ring theory in action

Let (\mathbb{E}, σ) be a hypergeometric *APS*-extension of (\mathbb{G}, σ) with ev : $\mathbb{E} \times \mathbb{N} \to \mathbb{K}$ and let $\tau : \mathbb{E} \to \mathbb{K}^{\mathbb{N}} / \sim$ be the \mathbb{K} -homomorphism given by

 $\tau(f) = (\operatorname{ev}(f, n))_{n \ge 0}.$

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 $\tau(f) = (\operatorname{ev}(f, n))_{n \ge 0}.$

Lemma

Let $W = \{T_1, \ldots, T_e\} \in \Sigma \Pi(\mathbb{G})$ with $T_i = expr(t_i)$. Then:

W is canonical reduced $\Leftrightarrow \tau$ is injective.
Difference ring theory in action

Let (\mathbb{E}, σ) be a hypergeometric APS-extension of (\mathbb{G}, σ) with $ev : \mathbb{E} \times \mathbb{N} \to \mathbb{K}$ and let $\tau : \mathbb{E} \to \mathbb{K}^{\mathbb{N}} / \sim$ be the \mathbb{K} -homomorphism given by

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Let
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W is canonical reduced $\Leftrightarrow \tau$ is injective.

Using difference ring theory we get the following crucial property: Theorem

$$\tau$$
 is injective $\Leftrightarrow \operatorname{const}_{\sigma} \mathbb{E} = \mathbb{K}.$

Example

For our difference field $\mathbb{G} = \mathbb{K}(x)$ with $\sigma(x) = x + 1$ and $\text{const}_{\sigma}\mathbb{K} = \mathbb{K}$ we have $\text{const}_{\sigma}\mathbb{K}(x) = \mathbb{K}$.

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Definition

A hypergeometric *APS*-extension (\mathbb{E}, σ) of (\mathbb{G}, σ) is called **hypergeometric** $R\Pi\Sigma$ -extension if

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Definition

A hypergeometric *APS*-extension (\mathbb{E}, σ) of (\mathbb{G}, σ) is called **hypergeometric** $R\Pi\Sigma$ -extension if

$$\operatorname{const}_{\sigma}\mathbb{E} = \mathbb{K}.$$

Theorem

Let $W = \{T_1, \ldots, T_e\} \subset \Sigma\Pi(\mathbb{G})$ be in sum-product reduced representation and shift-stable, i.e., for each $1 \leq i \leq e$ the arising sums and products in T_i are contained in $\{T_1, \ldots, T_{i-1}\}$. Then the following is equivalent:

- 1. There is a hypergeometric $R\Pi\Sigma$ -extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ equipped with an evaluation function ev with $T_i = \exp(t_i) \in \Sigma\Pi(\mathbb{G})$ for $1 \le i \le e$.
- 2. W is σ -reduced over \mathbb{G} .

This yields a strategy (actually the only strategy for shift-stable sets).

A Strategy to solve Problem SigmaReduce

Given: $A_1, \ldots, A_u \in \mathsf{SumProd}(\mathbb{G})$ with $\mathbb{G} = \mathbb{K}(x)$.

- Find: a σ -reduced set $W = \{T_1, \ldots, T_e\} \subset \Sigma \Pi(\mathbb{G})$ with $B_1 \ldots, B_u \in$ SumProd (W, \mathbb{G}) and $\delta_1, \ldots, \delta_u \in \mathbb{N}$ such that $A_i(n) = B_i(n)$ holds for all $n \geq \delta_i$ and $1 \leq i \leq r$.
 - 1. Construct $R\Pi\Sigma$ -extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ equipped with $ev : \mathbb{E} \times \mathbb{N} \to \mathbb{K}$ such that we get $a_1, \dots, a_u \in \mathbb{E}$ and $\delta_1, \dots, \delta_u \in \mathbb{N}$ with

$$A_i(n) = \operatorname{ev}(a_i, n) \quad \forall n \ge \delta_i.$$
(12)

2. Set $W = \{T_1, \ldots, T_e\}$ with $T_i := \exp(t_i) \in \Sigma \Pi(\mathbb{G})$ for $1 \le i \le e$.

- 3. Set $B_i := \exp(a_i) \in \operatorname{SumProd}(W, \mathbb{G})$ for $1 \le i \le u$.
- 4. Return W, (B_1, \ldots, B_u) and $(\delta_1, \ldots, \delta_u)$.

General picture:



General picture:

Part 1: Symbolic summation (a short introduction)

Part 2: Modeling of sequences with a term algebra (user interface)

Part 3: Modeling of sequences in difference rings (computer algebra)

Part 4: Construction of appropriate difference rings (advanced CA)

Part 5: Applications

A hypergeometric $APS\mbox{-extension}$ of $(\mathbb{K}(x),\sigma)$ is

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2][s_3] \dots$$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\begin{split} \sigma(x) &= x + 1 \\ \text{hypergeometric} &\leftrightarrow \sigma(p_1) = a_1 p_1 & a_1 \in \mathbb{K}(x)^* \\ \text{products} &\sigma(p_2) = a_2 p_2 & a_2 \in \mathbb{K}(x)^* \\ \vdots \\ \sigma(p_e) &= a_e p_e & a_e \in \mathbb{K}(x)^* \\ \text{oot of unity} &\gamma^{\mathbf{k}} &\leftrightarrow \sigma(\mathbf{z}) = \gamma \mathbf{z} & \mathbf{z}^{\lambda} = \mathbf{1} \\ \text{(nested) sum} &\leftrightarrow \sigma(s_1) = s_1 + f_1 & f_1 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z] \\ \sigma(s_2) &= s_2 + f_2 & f_2 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1] \\ \sigma(s_3) &= s_3 + f_3 & f_3 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2] \\ \vdots \\ \end{split}$$

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 \blacktriangleright Let (\mathbb{A},σ) be a difference ring with constant set

$$\operatorname{const}_{\sigma} \mathbb{A} := \{ k \in \mathbb{A} | \sigma(k) = k \}.$$

Note 1: $const_{\sigma}\mathbb{A}$ is a ring that contains \mathbb{Q}

Note 2: We always take care that $const_{\sigma}\mathbb{A}$ is a field

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► Adjoin a new variable t to A (i.e., A[t] is a polynomial ring).

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$$\sigma(t) = t + f$$
 for some $f \in \mathbb{A}$.

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Such a difference ring extension $(\mathbb{A}[t], \sigma)$ of (\mathbb{A}, σ) is called Σ^* -extension

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There are 2 cases:

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$$\exists g \in \mathbb{A} : \sigma(g) = g + f : (\mathbb{A}[t], \sigma) \text{ is a } \Sigma^*\text{-extension of } (\mathbb{A}, \sigma)$$

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2. $\exists g \in \mathbb{A} : \sigma(g) = g + f$: No need for a Σ^* -extension!

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Such a difference ring extension $(\mathbb{A}[t,\frac{1}{t}],\sigma)$ of (\mathbb{A},σ) is called Π -extension

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There are 3 cases:

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$$\exists g \in \mathbb{A} \setminus \{0\} \exists n \in \mathbb{Z} \setminus \{0\} : \sigma(g) = a^n g$$
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There are 3 cases:

- 1. $\nexists g \in \mathbb{A} \setminus \{0\} \nexists n \in \mathbb{Z} \setminus \{0\} : \sigma(g) = a^n g$: $(\mathbb{A}[t, \frac{1}{t}]), \sigma)$ is a Π -ext. of (\mathbb{A}, σ)
- 2. $\exists g \in \mathbb{A} \setminus \{0\} : \sigma(g) = a g :$ No need for a Π -extension!

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- 2. $\exists g \in \mathbb{A} \setminus \{0\} : \sigma(g) = a g \models$ No need for a Π -extension!
- 3. $\exists g \in \mathbb{A} \setminus \{0\} : \sigma(g) = a^n g \text{ only for } n \in \mathbb{Z} \setminus \{0, 1\} \models \bigcirc$

- ► Take the difference field $(\mathbb{K}(x), \sigma)$ with $\sigma|_{\mathbb{K}} = \text{id}$ and $\sigma(x) = x + 1$.
- Let $\alpha_1, \ldots, \alpha_r \in \mathbb{K}(x)^*$

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 \mathbb{E}

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$$\mathbb{E} = \mathbb{K}(x) \underbrace{[t_1, t_1^{-1}] \dots [t_e, t_e^{-1}]}_{\text{tower of Π-ext.}} \underbrace{[z]}_{(-1)^k \text{ or } \gamma^k}$$

with

▶
$$\frac{\sigma(t_i)}{t_i} \in \mathbb{K}(x)^*$$
 for $1 \le i \le e$
▶ $\sigma(z) = \gamma z$ and $z^{\lambda} = 1$ for some primitive λ th root of unity $\gamma \in \mathbb{K}^*$
▶ $\text{const}_{\sigma}\mathbb{E} = \mathbb{K}$

such that for $1 \leq i \leq r$ there are $g_i \in \mathbb{E}^*$ with

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with

▶
$$\frac{\sigma(t_i)}{t_i} \in \mathbb{K}(x)^*$$
 for $1 \le i \le e$
▶ $\sigma(z) = \gamma z$ and $z^{\lambda} = 1$ for some primitive λ th root of unity $\gamma \in \mathbb{K}^*$
▶ $\text{const}_{\sigma}\mathbb{E} = \mathbb{K}$

such that for $1 \le i \le r$ there are $g_i \in \mathbb{E}^*$ with $\boxed{\sigma(g_i) = \alpha_i \, g_i}$

Note: There are similar results for the q-rational, multi-basic and mixed case

A hypergeometric $R\Pi\Sigma$ -extension of $(\mathbb{K}(x), \sigma)$ is

a ring

$$\mathbb{A} := \mathbb{K}(x)[p_1, p_1^{-1}][p_2, p_2^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2][s_3] \dots$$

 \blacktriangleright with an automorphism where $\sigma(c)=c$ for all $c\in\mathbb{K}$ and where

$$\begin{split} \sigma(x) &= x+1 \\ \text{hypergeometric} &\leftrightarrow \sigma(p_1) = a_1 \, p_1 & a_1 \in \mathbb{K}(x)^* \\ \text{products} &\sigma(p_2) = a_2 p_2 & a_2 \in \mathbb{K}(x)^* \\ \vdots \\ \sigma(p_e) &= a_e p_e & a_e \in \mathbb{K}(x)^* \\ \gamma \text{ is a primitive } \lambda \text{th} \quad \gamma^{\mathbf{k}} &\leftrightarrow \sigma(\mathbf{z}) = \gamma \, \mathbf{z} & \mathbf{z}^{\lambda} = \mathbf{1} \\ \text{(nested) sum} &\leftrightarrow \sigma(s_1) = s_1 + f_1 \quad f_1 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z] \\ \sigma(s_2) &= s_2 + f_2 \quad f_2 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1] \\ \sigma(s_3) &= s_3 + f_3 \quad f_3 \in \mathbb{K}(x)[p_1, p_1^{-1}] \dots [p_e, p_e^{-1}][z][s_1][s_2] \\ \vdots \\ \text{such that const}_{\sigma} \mathbb{E} = \mathbb{K} \end{split}$$

This yields a strategy (actually the only strategy for shift-stable sets).

A Strategy to solve Problem SigmaReduce

- Given: $A_1, \ldots, A_u \in \mathsf{SumProd}(\mathbb{G})$ with $\mathbb{G} = \mathbb{K}(x)$.
- Find: a σ -reduced set $W = \{T_1, \ldots, T_e\} \subset \Sigma \Pi(\mathbb{G})$ with $B_1 \ldots, B_u \in SumProd(W, \mathbb{G})$ and $\delta_1, \ldots, \delta_u \in \mathbb{N}$ such that $A_i(n) = B_i(n)$ holds for all $n \geq \delta_i$ and $1 \leq i \leq r$.
 - 1. Construct $R\Pi\Sigma$ -extension (\mathbb{E}, σ) of (\mathbb{G}, σ) with $\mathbb{E} = \mathbb{G}\langle t_1 \rangle \dots \langle t_e \rangle$ equipped with $ev : \mathbb{E} \times \mathbb{N} \to \mathbb{K}$ such that we get $a_1, \dots, a_u \in \mathbb{E}$ and $\delta_1, \dots, \delta_u \in \mathbb{N}$ with

$$A_i(n) = \operatorname{ev}(a_i, n) \quad \forall n \ge \delta_i.$$
(12)

- 2. Set $W = \{T_1, \ldots, T_e\}$ with $T_i := \exp(t_i) \in \Sigma \Pi(\mathbb{G})$ for $1 \le i \le e$.
- 3. Set $B_i := \exp(a_i) \in \operatorname{SumProd}(W, \mathbb{G})$ for $1 \le i \le u$.
- 4. Return W, (B_1, \ldots, B_u) and $(\delta_1, \ldots, \delta_u)$.

This yields a strategy (actually the only strategy for shift-stable sets).

An Algorithm to solve Problem SigmaReduce

- Given: $A_1, \ldots, A_u \in \mathsf{SumProd}(\mathbb{G})$ with $\mathbb{G} = \mathbb{K}(x)$.
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Telescoping

GIVEN
$$f(k) = S_1(k)$$
.
FIND $g(k)$:

$$f(k) = g(k+1) - g(k)$$
for all $1 \le k \le n$ and $n \ge 0$.

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Sigma computes

$$g(k) = (S_1(k) - 1)k.$$
Telescoping

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FIND $g(k)$:

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for all $1 \le k \le n$ and $n \ge 0$.

Summing this equation over $k \mbox{ from } 1 \mbox{ to } n \mbox{ gives }$

$$\boxed{\sum_{k=1}^{n} S_1(k)} = \boxed{g(n+1) - g(1)} = (S_1(n+1) - 1)(n+1).$$

FIND a closed form for

$$\sum_{k=1}^{n} S_1(k).$$

A difference ring for the summand

Consider a ring

A

FIND a closed form for

$$\sum_{k=1}^{n} S_1(k).$$

A difference ring for the summand

Consider a ring

$$\mathbb{A}:=\mathbb{Q}$$

$$\sigma(c) = c \quad \forall c \in \mathbb{Q},$$

FIND a closed form for

$$\sum_{k=1}^n S_1(\boldsymbol{k}).$$

A difference ring for the summand

Consider a ring

$$\mathbb{A} := \mathbb{Q}(x)$$

$$\begin{aligned} \sigma(c) &= c \quad \forall c \in \mathbb{Q}, \\ \sigma(x) &= x+1, \end{aligned} \qquad \qquad \qquad \mathcal{S} \, k = k+1, \end{aligned}$$

FIND a closed form for

$$\sum_{k=1}^n S_1(k).$$

A difference ring for the summand

Consider a ring

$$\mathbb{A} := \mathbb{Q}(x)[h]$$

$$\sigma(c) = c \quad \forall c \in \mathbb{Q},$$

$$\sigma(x) = x + 1, \qquad \qquad \mathcal{S} k = k + 1,$$

$$\sigma(h) = h + \frac{1}{x + 1}, \qquad \qquad \mathcal{S} S_1(k) = S_1(k) + \frac{1}{k + 1}.$$

FIND $g \in \mathbb{A}$:

$$\sigma(g) - g = h.$$





This gives $\label{eq:gk} \boxed{g(k+1)-g(k)=S_1(k)}$ with $g(k)=(S_1(k)-1)k.$





Hence,

$$(S_1(n+1)-1)(n+1) = \sum_{k=1}^n S_1(k).$$

$$\sigma(g) - g = h.$$

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Degree bound: COMPUTE $b \ge 0$:

$$\forall g \in \mathbb{Q}(x)[h] \quad \sigma(g) - g = h \quad \Rightarrow \quad \deg(g) \le b.$$

$$\sigma(g) - g = h.$$



$$\sigma(g) - g = h.$$

Degree bound: COMPUTE $b \ge 0$:

b=2

$$\forall g \in \mathbb{Q}(x)[h] \quad \sigma(g) - g = h \quad \Rightarrow \quad \deg(g) \le b.$$

Polynomial Solution: FIND

$$g = g_2 h^2 + g_1 h + g_0 \in \mathbb{Q}(x)[h].$$

$$\sigma(g) - g = h$$

$$\begin{bmatrix} \sigma(g_2 h^2 + g_1 h + g_0) \end{bmatrix} - \begin{bmatrix} g_2 h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$\begin{bmatrix} \sigma(g_2 h^2) + \sigma(g_1 h + g_0) \\ - [g_2 h^2 + g_1 h + g_0] = h \end{bmatrix}$$

$$\begin{bmatrix} \sigma(g_2) \, \sigma(h^2) + \sigma(g_1 h + g_0) \\ - \begin{bmatrix} g_2 \, h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$\begin{bmatrix} \sigma(g_2) \, \sigma(h)^2 + \sigma(g_1 h + g_0) \\ - \begin{bmatrix} g_2 \, h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$\begin{bmatrix} \sigma(g_2) \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1 h + g_0) \\ - [g_2 h^2 + g_1 h + g_0] = h \end{bmatrix}$$

$$\begin{bmatrix} \sigma(g_2) \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1 h + g_0) \end{bmatrix} = h$$

$$- \begin{bmatrix} g_2 h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$\sigma(g_2) - g_2 = 0$$

$$\left[\sigma(g_2) \left(h + \frac{1}{x+1} \right)^2 + \sigma(g_1 h + g_0) \right]$$

- $\left[g_2 h^2 + g_1 h + g_0 \right] = h$
 $g_2 = c \in \mathbb{Q}$

$$\begin{bmatrix} \sigma(g_2) \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1 h + g_0) \end{bmatrix} - \begin{bmatrix} g_2 h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$[\sigma(g_2) - g_2 = 0]$$

$$\begin{bmatrix} \sigma(g_2) - g_2 = 0 \end{bmatrix}$$

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$$\begin{bmatrix} \sigma(c) \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1 h + g_0) \end{bmatrix}$$

$$- \begin{bmatrix} c h^2 + g_1 h + g_0 \end{bmatrix} = h$$

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$$- \begin{bmatrix} g_2 h^2 + g_1 h + g_0 \end{bmatrix} = h$$

$$\begin{bmatrix} \sigma(g_2) - g_2 = 0 \end{bmatrix}$$

$$g_2 = c \in \mathbb{Q}$$

$$\begin{bmatrix} c \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1h + g_0) \end{bmatrix}$$

$$- \begin{bmatrix} c h^2 + g_1 h + g_0 \end{bmatrix} = h$$

ANSATZ
$$g = g_2 h^2 + g_1 h + g_0 \in \mathbb{Q}(x)[h]$$

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$$\begin{bmatrix} \sigma(g_2) \left(h + \frac{1}{x+1}\right)^2 + \sigma(g_1 h + g_0) \end{bmatrix} = h$$
 coeff. comp.

$$\sigma(g_1 h + g_0) - (g_1 h + g_0) = h - c \begin{bmatrix} \frac{2h(x+1)+1}{(x+1)^2} \end{bmatrix}$$
 coeff. comp.

$$\sigma(g_1) - g_1 = 1 - c \frac{2}{x+1}$$

ANSATZ
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$$\begin{split} \left[\sigma(g_2)\left(h+\frac{1}{x+1}\right)^2 + \sigma(g_1h+g_0)\right] & -\left[g_2\,h^2 + g_1h + g_0\right] = h & \text{coeff. comp.} \\ & \sigma(g_2) - g_2 = 0 \\ g_2 = c \in \mathbb{Q} \\ \sigma(g_1\,h+g_0) - (g_1\,h+g_0) = h - c\left[\frac{2h(x+1)+1}{(x+1)^2}\right] & \text{coeff. comp.} \\ & \sigma(g_1) - g_1 = 1 - c\,\frac{2}{x+1} \\ & \sigma(g_1) - g_1 = 1 - c\,\frac{2}{x+1} \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) - g_1 = x + d \\ & \sigma(g_1) -$$

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$$g = g_2 h^2 + g_1 h + g_0 \in \mathbb{Q}(x)[h]$$

$$\begin{split} \left[\sigma(g_2)\left(h+\frac{1}{x+1}\right)^2 + \sigma(g_1h+g_0)\right] & -\left[g_2h^2 + g_1h + g_0\right] = h & \text{coeff. comp.} \\ & \sigma(g_2) - g_2 = 0 \\ g_2 = c \in \mathbb{Q} \\ \sigma(g_1h+g_0) - (g_1h+g_0) = h - c\left[\frac{2h(x+1)+1}{(x+1)^2}\right] & \text{coeff. comp.} \\ & \sigma(g_1) - g_1 = 1 - c\frac{2}{x+1} \\ & \sigma(g_0) - g_0 = -1 - d\frac{1}{x+1} & c = 0, \quad g_1 = x+d \\ & d \in \mathbb{Q} \\ \end{split}$$

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Hence,

$$(S_1(n+1)-1)(n+1) = \sum_{k=1}^n S_1(k).$$

▶ the mixed multibasic hypergeometric case: $\mathbb{G} := \mathbb{K}(x, x_1, \dots, x_v)$ with $\mathbb{K} = K(q_1, \dots, q_v)$ For $f = \frac{p}{q} \in \mathbb{G}$ with $p, q \in \mathbb{K}[x, x_1, \dots, x_v]$ where $q \neq 0$ and p, q being coprime we define

$$\operatorname{ev}(f,k) = \begin{cases} 0 & \text{if } q(k,q_1^k,\dots,q_v^k) = 0\\ \frac{p(k,q_1^k,\dots,q_v^k)}{q(k,q_1^k,\dots,q_v^k)} & \text{if } q(k,q_1^k,\dots,q_v^k) \neq 0 \end{cases}$$

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- ▶ simple products: $\operatorname{Prod}^*(\mathbb{G})$ is the smallest set that contains 1 with: 1. If $r \in \mathcal{R}$ then $\operatorname{RPow}(r) \in \operatorname{Prod}^*(\mathbb{G})$.
- 2. If $f \in \mathbb{G}^*$, $l \in \mathbb{N}$ with $l \ge Z(f)$ then $\operatorname{Prod}(l, f) \in \operatorname{Prod}^*(\mathbb{G})$.
- 3. If $p, q \in \mathsf{Prod}^*(\mathbb{G})$ then $p \odot q \in \mathsf{Prod}^*(\mathbb{G})$.
- 4. If $p \in \operatorname{Prod}^*(\mathbb{G})$ and $z \in \mathbb{Z} \setminus \{0\}$ then $p^{\bigcirc}z \in \operatorname{Prod}^*(\mathbb{G})$.

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- 4. If $p \in \operatorname{Prod}^*(\mathbb{G})$ and $z \in \mathbb{Z} \setminus \{0\}$ then $p^{\bigcirc}z \in \operatorname{Prod}^*(\mathbb{G})$.

For further details see

Term Algebras, Canonical Representations and Difference Ring Theory for Symbolic Summation. To appear in: Anti-Differentiation and the Calculation of Feynman Amplitudes, J. Blümlein and C. Schneider (ed.), Texts and Monographs in Symbolic Computation, 2021. Springer, arXiv:2102.01471 [cs.SC]

General picture:

Part 1: Symbolic summation (a short introduction)

Part 2: Modeling of sequences with a term algebra (user interface)

Part 3: Modeling of sequences in difference rings (computer algebra)

Part 4: Construction of appropriate difference rings (advanced CA)

Part 5: Applications



$$S_1(n) = \sum_{i=1}^n \frac{1}{i} \quad (=H_n)$$

Arose in the context of

I. Bierenbaum, J. Blümlein, and S. Klein, **Evaluating two-loop massive operator matrix** elements with Mellin-Barnes integrals. 2006

A warm-up example: simplify

$$\sum_{k=0}^{\infty} \sum_{j=0}^{\infty} \left(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)!(j+n+1)!(k+n+1)!} + \frac{j!k!(j+k+n)!(-S_1(j)+S_1(j+k)+S_1(j+n)-S_1(j+k+n))}{(j+k+1)!(j+n+1)!(k+n+1)!} \right)$$
FIND $g(j)$:

$$f(j) = g(j+1) - g(j)$$


FIND g(j):

$$f(j) = g(j+1) - g(j)$$

↑ summation package Sigma

$$g(j) = \frac{(j+k+1)(j+n+1)j!k!(j+k+n)!\left(S_1(j)-S_1(j+k)-S_1(j+n)+S_1(j+k+n)\right)}{kn(j+k+1)!(j+n+1)!(k+n+1)!}$$

A warm-up example: simplify

$$\sum_{k=0}^{\infty} \sum_{j=0}^{\infty} \left(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)(j+k+1)!(j+n+1)!(k+n+1)!} + \frac{j!k!(j+k+n)!(-S_1(j)+S_1(j+k)+S_1(j+n)-S_1(j+k+n))}{(j+k+1)!(j+n+1)!(k+n+1)!} \right)$$
FIND: (i)

FIND g(j):

$$f(j) = g(j+1) - g(j)$$

Summing the telescoping equation over \boldsymbol{j} from $\boldsymbol{0}$ to \boldsymbol{a} gives

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$$f(j) = g(j+1) - g(j)$$

Summing the telescoping equation over \boldsymbol{j} from $\boldsymbol{0}$ to \boldsymbol{a} gives

$$\begin{split} &\sum_{j=0}^{a} f(j) = g(a+1) - g(0) \\ &= \underbrace{\frac{(a+1)!(k-1)!(a+k+n+1)!(S_1(a) - S_1(a+k) - S_1(a+n) + S_1(a+k+n))}{n(a+k+1)!(a+n+1)!(k+n+1)!}}_{\substack{+ \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)n!} + \frac{(2a+k+n+2)a!k!(a+k+n)!}{(a+k+1)(a+n+1)(a+k+1)!(a+n+1)!(k+n+1)!}}_{a \to \infty} \end{split}$$



$$\sum_{j=0}^{\infty} f(j) = \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}$$

ln[8] := << Sigma.m

Sigma - A summation package by Carsten Schneider \bigodot RISC-Linz

$$\begin{split} & \ln[9]:= mySum = \sum_{j=0}^{a} \Big(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)(j+k+1)!(j+n+1)!(k+n+1)!} + \\ & \frac{j!k!(j+k+n)!\left(-S_{1}[j]+S_{1}[j+k]+S_{1}[j+n]-S_{1}[j+k+n]\right)}{(j+k+1)!(j+n+1)!(k+n+1)!} \Big); \end{split}$$

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Sigma - A summation package by Carsten Schneider © RISC-Linz

$$\begin{split} & \ln[9]:= mySum = \sum_{j=0}^{a} \Big(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)(j+k+1)!(j+n+1)!(k+n+1)!} + \\ & \frac{j!k!(j+k+n)!\left(-S_{1}[j]+S_{1}[j+k]+S_{1}[j+n]-S_{1}[j+k+n]\right)}{(j+k+1)!(j+n+1)!(k+n+1)!} \Big); \end{split}$$

$$\begin{split} & \text{In[10]:= res = SigmaReduce[mySum]} \\ & \text{Out[10]=} \quad \frac{(a+1)!(k-1)!(a+k+n+1)!\left(S_1[a]-S_1[a+k]-S_1[a+n]+S_1[a+k+n]\right)}{n(a+k+1)!(a+n+1)!(k+n+1)!} + \\ & \frac{S_1[k]+S_1[n]-S_1[k+n]}{kn(k+n+1)n!} + \frac{(2a+k+n+2)a!k!(a+k+n)!}{(a+k+1)(a+n+1)!(a+n+1)!(k+n+1)!} \end{split}$$

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Sigma - A summation package by Carsten Schneider \bigodot RISC-Linz

$$\begin{split} & \ln[9]:= mySum = \sum_{j=0}^{a} \Big(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)(j+k+1)!(j+n+1)!(k+n+1)!} + \\ & \frac{j!k!(j+k+n)!\left(-S_{1}[j]+S_{1}[j+k]+S_{1}[j+n]-S_{1}[j+k+n]\right)}{(j+k+1)!(j+n+1)!(k+n+1)!} \Big); \end{split}$$

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 ${\tt ln[11]:= SigmaLimit[res, \{n\}, a]}$

$$\label{eq:out_11} \text{Out}[11] = \quad \frac{1}{n!} \frac{S_1[k] + S_1[n] - S_1[k+n]}{kn(k+n+1)}$$



$$\sum_{j=0}^{\infty} f(j) = \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}$$



$$\sum_{k=1}^{\infty} \sum_{j=0}^{\infty} f(j) = \frac{1}{n!} \sum_{k=1}^{\infty} \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}$$

Telescoping

GIVEN

$$\mathsf{A}(n) := \sum_{k=1}^{a} \underbrace{\frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}}_{=: f(k)}.$$

FIND g(k):

$$\boxed{g(k+1) - g(k)} = \boxed{f(k)}$$

for all $0 \le k \le n$ and all $n \ge 0$.

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Zeilberger's creative telescoping paradigm GIVEN $A(n) := \sum_{k=1}^{a} \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}.$

FIND
$$g(n,k)$$

$$\boxed{g(n,k+1) - g(n,k)} = \boxed{f(n,k)}$$

for all $0 \le k \le n$ and all $n \ge 0$.



=: f(n,k)

A(n) :=
$$\sum_{k=1}^{a} \underbrace{\frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}}_{=:f(n,k)}$$
.

FIND g(n,k) and $c_0(n), c_1(n)$:

$$\boxed{g(n,k+1) - g(n,k)} = \boxed{c_0(n)f(n,k) + c_1(n)f(n+1,k)}$$

for all $0 \le k \le n$ and all $n \ge 0$.

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for all $0 \le k \le n$ and all $n \ge 0$.

Sigma computes: $c_0(n) = -n$, $c_1(n) = (n+2)$ and

$$g(n,k) = \frac{kS_1(k) + (-n-1)S_1(n) - kS_1(k+n) - 2}{(k+n+1)(n+1)^2}$$

GIVEN

$$A(n) := \sum_{k=1}^{a} \underbrace{\frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}}_{=:f(n,k)}$$

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$$g(n,k+1) - g(n,k) = c_0(n)f(n,k) + c_1(n)f(n+1,k)$$

for all $0 \le k \le n$ and all $n \ge 0$.

$$\boxed{g(n, a+1) - g(n, 1)} = \sum_{k=1}^{a} \left[c_0(n) f(n, k) + c_1(n) f(n+1, k) \right]$$

GIVEN

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$$g(n, a+1) - g(n, 1) = c_0(n) \sum_{k=1}^{a} f(n, k) + c_1(n) \sum_{k=1}^{a} f(n+1, k)$$

GIVEN

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for all $0 \le k \le n$ and all $n \ge 0$.

$$g(n, a+1) - g(n, 1) = c_0(n) \mathsf{A}(n) + c_1(n) \mathsf{A}(n+1)$$

GIVEN

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for all $0 \le k \le n$ and all $n \ge 0$.

$$\boxed{\begin{array}{c}g(n, a+1) - g(n, 1) \\ || \\ (a+1)(S_1(a) + S_1(n) - S_1(a+n)) \\ (n+1)^2(a+n+2) \\ + \frac{a(a+1)}{(n+1)^3(a+n+1)(a+n+2)} \end{array}} - n\mathsf{A}(n) + (2+n)\mathsf{A}(n+1)$$



$$(n+2)\mathbf{A}(n+1) - n\mathbf{A}(n) = \frac{(n+1)S_1(n) + 1}{(n+1)^3}$$

recurrence solver
$$A(n) = \sum_{k=1}^{\infty} \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)} \in \frac{\{c \times \frac{1}{n(n+1)} + \frac{S_1(n)^2 + S_2(n)}{2n(n+1)} | c \in \mathbb{R}\}$$
where

$$S_1(n) = \sum_{i=1}^n \frac{1}{i} \qquad S_2(n)$$

$$S_2(n) = \sum_{i=1}^n \frac{1}{i^2}$$

Г

$$(n+2)\mathbf{A}(n+1) - n\mathbf{A}(n) = \frac{(n+1)S_1(n) + 1}{(n+1)^3}$$
Summation package Sigma
(based on difference field/ring algorithms/theory
see, e.g., Abramov, Karr 1981, Bronstein 2000, Schneider 2001/2004/2005a-c/2007/2008/201Da-c)

$$\overline{A(n) = \sum_{k=1}^{\infty} \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}} = \underbrace{\begin{array}{c} 0 \times \frac{1}{n(n+1)} \\ + \frac{S_1(n)^2 + S_2(n)}{2n(n+1)} \\ + \frac{S_1(n)^2 + S_2(n)}{2n(n+1)} \end{array}}_{\text{where}}$$
where

$$S_1(n) = \sum_{i=1}^{n} \frac{1}{i} \qquad S_2(n) = \sum_{i=1}^{n} \frac{1}{i^2}$$

$$\ln[12] = mySum = \sum_{k=1}^{n} \frac{S[1,k] + S[1,n] - S[1,k+n]}{kn(k+n+1)};$$

$$\ln[12] = \mathbf{mySum} = \sum_{k=1}^{a} \frac{\mathbf{S}[1,k] + \mathbf{S}[1,n] - \mathbf{S}[1,k+n]}{\mathbf{kn}(\mathbf{k}+\mathbf{n}+1)};$$

ln[13]:= rec = GenerateRecurrence[mySum, n][[1]]

$$\mathsf{Out}[13] = n\mathsf{SUM}[n] + (1+n)(2+n)\mathsf{SUM}[n+1] = = \frac{(a+1)(S[1,a]+S[1,n]-S[1,a+n])}{(n+1)^2(a+n+2)n!} + \frac{a(a+1)(A+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2)(a+n+2$$

$$\ln[12] = \mathbf{mySum} = \sum_{k=1}^{a} \frac{\mathbf{S}[1,k] + \mathbf{S}[1,n] - \mathbf{S}[1,k+n]}{kn(k+n+1)};$$

ln[13]:=rec = GenerateRecurrence[mySum, n][[1]]

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 $ln[14] = rec = LimitRec[rec, SUM[n], \{n\}, a]$

 $\label{eq:out_state} Out[14]= -n\mathsf{SUM}[n] + (1+n)(2+n)\mathsf{SUM}[n+1] = = \frac{(n+1)\mathsf{S}[1,n]+1}{(n+1)^3}$

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Solve a recurrence

$$\begin{split} & \text{In[15]:= recSol = SolveRecurrence[rec, SUM[n]]} \\ & \text{Out[15]= } \{\{0, \frac{1}{n(n+1)}\}, \{1, \frac{S[1,n]^2 + \sum_{i=1}^{n} \frac{1}{i^2}}{2n(n+1)}\}\} \end{split}$$

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$$\begin{split} & \text{In[15]:= recSol = SolveRecurrence[rec, SUM[n]]} \\ & \text{Out[15]= } \{\{0, \frac{1}{n(n+1)}\}, \{1, \frac{S[1,n]^2 + \sum_{i=1}^{n} \frac{1}{i^2}}{2n(n+1)}\}\} \end{split}$$

Combine the solutions

 $\label{eq:ln[16]:=FindLinearCombination[recSol, {1, {1/2}, n, 2]} \\ \texttt{Out[16]=} \quad \frac{\texttt{S}[1, n]^2 + \sum_{i=1}^n \frac{1}{i^2}}{2n(n+1)}$



$$\sum_{k=1}^{\infty} \sum_{j=0}^{\infty} f(j) = \frac{1}{n!} \sum_{k=1}^{\infty} \frac{S_1(k) + S_1(n) - S_1(k+n)}{kn(k+n+1)}$$
$$= \frac{1}{n!} \frac{S_1(n)^2 + S_2(n)}{2n(n+1)}$$

where

$$S_1(n) = \sum_{i=1}^n \frac{1}{i}$$
 $S_2(n) = \sum_{i=1}^n \frac{1}{i^2}$

A warm-up example: simplify

$$\sum_{k=0}^{\infty} \sum_{j=0}^{\infty} \left(\frac{(2j+k+n+2)j!k!(j+k+n)!}{(j+k+1)(j+n+1)!(j+n+1)!(k+n+1)!} + \frac{j!k!(j+k+n)!(-S_1(j)+S_1(j+k)+S_1(j+n)-S_1(j+k+n))}{(j+k+1)!(j+n+1)!(k+n+1)!} \right)$$

$$\sum_{k=0}^{\infty} \sum_{j=0}^{\infty} f(n,k,j) = \frac{S_1(n)^2 + 3S_2(n)}{2n(n+1)!}$$

where

$$S_1(n) = \sum_{i=1}^n \frac{1}{i}$$
 $S_2(n) = \sum_{i=1}^n \frac{1}{i^2}$

1. Creative telescoping (for the special case of hypergeometric terms see Zeilberger's algorithm (1991)) GIVEN a definite sum

$$A(n) = \sum_{k=0}^{n} f(n,k);$$

 $f(n,k){:}$ indefinite nested product-sum in $k{;}$ $n{:}$ extra parameter

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2. <u>Recurrence solving</u>

GIVEN a recurrence

 $a_0(n), \ldots, a_d(n), h(n)$: indefinite nested product-sum expressions.

$$a_0(n)A(n) + \dots + a_d(n)A(n+d) = h(n);$$

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FIND all solutions expressible by indefinite nested products/sums (Abramov/Bronstein/Petkovšek/CS, in preparation)

3. Find a "closed form"

A(n)=combined solutions in terms of indefinite nested sums.

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} \binom{j+1}{r} \binom{-j+n+r-2}{s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!}$$

Simple sum

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!}$$

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \left| \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!} \right|$$

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!}$$

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \left| \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!} \right| \\ \left| \left(\frac{j+1}{r} \right) \left(\frac{(-1)^r (-j+n-2)! r!}{(n+1)(-j+n+r-1)(-j+n+r)!} + \frac{(-1)^{n+r} (j+1)! (-j+n-2)! (-j+n-1)_r r!}{(n-1)n(n+1)(-j+n+r)! (-j-1)_r (2-n)_j} \right) \right|$$

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {j-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!} \\ = \frac{1}{||} \\ \sum_{j=0}^{n-2} \left[\sum_{r=0}^{j+1} {j+1 \choose r} \left(\frac{(-1)^r (-j+n-2)! r!}{(n+1)(-j+n+r-1)(-j+n+r)!} + \frac{(-1)^{n+r} (j+1)! (-j+n-2)! (-j+n-1)_r r!}{(n-1)n(n+1)(-j+n+r)! (-j-1)_r (2-n)_j} \right) \right]$$
Part 1: Crucial summation paradigms

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!} \\ = \sum_{j=0}^{n-2} \left[\sum_{r=0}^{j+1} {j+1 \choose r} \left(\frac{(-1)^r (-j+n-2)! r!}{(n+1)(-j+n+r-1)(-j+n+r)!} + \frac{(-1)^{n+r} (j+1)! (-j+n-2)! (-j+n-1)_r r!}{(n-1)n(n+1)(-j+n+r)! (-j-1)_r (2-n)_j} \right) \right]$$

$$\left[\frac{\left(\frac{n^2-n+1}{(n-1)^2n^2(n+1)(2-n)_j}+\frac{\sum\limits_{i=1}^{j}\frac{(2-n)_i}{(-i+n-1)^2(i+1)!}}{(n+1)(2-n)_j}+\frac{(-1)^{j+n}(-j-2)(-j+n-2)!}{(j-n+1)(n+1)^2n!}\right)(j+1)!-\frac{1}{(n+1)^2(-j+n-1)}\right]$$

Part 1: Crucial summation paradigms

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!}$$

$$\sum_{j=0}^{n-2} \left(\left(\frac{n^2 - n + 1}{(n-1)^2 n^2 (n+1)(2-n)_j} + \frac{\sum_{i=1}^j \frac{(2-n)_i}{(-i+n-1)^2 (i+1)!}}{(n+1)(2-n)_j} + \frac{(2-n)_i}{(n+1)(2-n)_j} + \frac{(2-n)_i}{(n+1)(2-n)_i} + \frac{(2-$$

$$\frac{(-1)^{j+n}(-j-2)(-j+n-2)!}{(j-n+1)(n+1)^2n!}\Big)(j+1)! - \frac{1}{(n+1)^2(-j+n-1)}\Big)$$

Part 1: Crucial summation paradigms

$$\sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {-j+n+r-2 \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!}$$

$$\sum_{j=0}^{n-2} \left(\left(\frac{n^2 - n + 1}{(n-1)^2 n^2 (n+1)(2-n)_j} + \frac{\sum_{i=1}^j \frac{(2-n)_i}{(-i+n-1)^2 (i+1)!}}{(n+1)(2-n)_j} + \right) \right)$$

$$\frac{(-1)^{j+n}(-j-2)(-j+n-2)!}{(j-n+1)(n+1)^2n!}\Big)(j+1)! - \frac{1}{(n+1)^2(-j+n-1)}\Big)$$

$$\frac{-n^2 - n - 1}{n^2(n+1)^3} + \frac{(-1)^n \left(n^2 + n + 1\right)}{n^2(n+1)^3} - \frac{2S_{-2}(n)}{n+1} + \frac{S_1(n)}{(n+1)^2} + \frac{S_2(n)}{-n-1}$$

Note: $S_a(n) = \sum_{i=1}^N \frac{\operatorname{sign}(a)^i}{i^{|a|}}$, $a \in \mathbb{Z} \setminus \{0\}$.

ln[1] = << Sigma.m

Sigma - A summation package by Carsten Schneider (C) RISC-Linz

In[2] := << HarmonicSums.m

HarmonicSums by Jakob Ablinger ⓒ RISC-Linz

$ln[3] := << {\bf EvaluateMultiSums.m}$

EvaluateMultiSums by Carsten Schneider ⓒ RISC-Linz

 $ln[1] = \langle \mathbf{Sigma.m} \rangle$

Sigma - A summation package by Carsten Schneider (C) RISC-Linz

In[2]:= << HarmonicSums.m

HarmonicSums by Jakob Ablinger ⓒ RISC-Linz

In[3] := << EvaluateMultiSums.m

EvaluateMultiSums by Carsten Schneider ⓒ RISC-Linz

$$\ln[4] = mySum = \sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} \binom{j+1}{r} \binom{-j+n+r-2}{s} (-j+n-2)!r!}{(n-s)(s+1)(-j+n+r)!};$$

 $\label{eq:ln[5]:=EvaluateMultiSum[mySum, \{\}, \{n\}, \{1\}]$

ln[1] = << Sigma.m

Sigma - A summation package by Carsten Schneider (C) RISC-Linz

In[2] = << HarmonicSums.m

HarmonicSums by Jakob Ablinger ⓒ RISC-Linz

In[3] := << EvaluateMultiSums.m

EvaluateMultiSums by Carsten Schneider ⓒ RISC-Linz

$$\ln[4] = mySum = \sum_{j=0}^{n-2} \sum_{r=0}^{j+1} \sum_{s=0}^{n-j+r-2} \frac{(-1)^{r+s} {j+1 \choose r} {(-j+n+r-2) \choose s} (-j+n-2)! r!}{(n-s)(s+1)(-j+n+r)!};$$

 $\label{eq:ln[5]:=EvaluateMultiSum[mySum, \{\}, \{n\}, \{1\}]$

$$\text{Out[5]}= \quad \frac{-n^2-n-1}{n^2(n+1)^3} + \frac{(-1)^n\left(n^2+n+1\right)}{n^2(n+1)^3} - \frac{2S[-2,n]}{n+1} + \frac{S[1,n]}{(n+1)^2} + \frac{S[2,n]}{-n-1} + \frac{S[2,n]}{2} + \frac{S[2,n$$

Application: The simplification of Feynman integrals



Behavior of particles



Behavior of particles

 $\rightarrow \int \Phi(N,\epsilon,x) dx$ Feynman integrals

$$\int_0^1 x \, dx = ?$$





$$\int_0^1 x^2 dx = ?$$





$$\int_0^1 x^3 dx = ?$$



$$\int_{0}^{1} x^{1} dx = \int_{0}^{1} x^{N} dx = \frac{1}{N+1}$$

$$\int_{0}^{1} x^{2} dx = \int_{0}^{1} x^{3} dx = \int_{0}^{1} x^{3}$$

0.4

0.6

0.8

1.0

$$\int_0^1 x^N \, dx$$

$$\int_0^1 x^N (1+x)^N \, dx$$

$$\int_0^1 \frac{x^N (1+x)^N}{(1-x)^{1+\varepsilon}} \, dx$$

$$\int_0^1 \int_0^1 \frac{x_1^N (1+x_1)^N}{(1-x_1)^{1+\varepsilon}} \dots \, dx_1 \, dx_2$$

$$\int_0^1 \int_0^1 \int_0^1 \frac{x_1^N (1+x_1)^N}{(1-x_1)^{1+\varepsilon}} \dots \, dx_1 \, dx_2 \, dx_3$$

$$\int_0^1 \int_0^1 \int_0^1 \int_0^1 \frac{x_1^N (1+x_1)^N}{(1-x_1)^{1+\varepsilon}} \dots \, dx_1 \, dx_2 \, dx_3 \, dx_4$$

$$\int_0^1 \int_0^1 \int_0^1 \int_0^1 \int_0^1 \frac{x_1^N (1+x_1)^N}{(1-x_1)^{1+\varepsilon}} \dots \, dx_1 \, dx_2 \, dx_3 \, dx_4 \, dx_5$$

$$\int_0^1 \int_0^1 \int_0^1 \int_0^1 \int_0^1 \int_0^1 \frac{x_1^N (1+x_1)^N}{(1-x_1)^{1+\varepsilon}} \dots \, dx_1 \, dx_2 \, dx_3 \, dx_4 \, dx_5 \, dx_6$$

$$\sum_{j=0}^{N-3} \sum_{k=0}^{j} \binom{N-1}{j+2} \binom{j+1}{k+1} \\ \times \int_{0}^{1} \int_{0}^{1} \int_{0}^{1} \int_{0}^{1} \int_{0}^{1} \int_{0}^{1} \frac{x_{1}^{N}(1+x_{1})^{N-j+k}}{(1-x_{1})^{1+\varepsilon}} \dots dx_{1} dx_{2} dx_{3} dx_{4} dx_{5} dx_{6}$$





Behavior of particles

 $\rightarrow \int \Phi(N,\epsilon,x) dx$ Feynman integrals





Example 1:

massive 3-loop ladder integrals





$$= F_{-3}(N)\varepsilon^{-3} + F_{-2}(N)\varepsilon^{-2} + F_{-1}(N)\varepsilon^{-1} + F_{0}(N)$$



$$= F_{-3}(N)\varepsilon^{-3} + F_{-2}(N)\varepsilon^{-2} + F_{-1}(N)\varepsilon^{-1} + F_{0}(N)$$

Simplify

$$\sum_{j=0}^{N-3} \sum_{k=0}^{j} \sum_{l=0}^{k} \sum_{q=0}^{-j+N-3} \sum_{s=1}^{-l+N-q-3} \sum_{r=0}^{-l+N-q-s-3} (-1)^{-j+k-l+N-q-3} \times \frac{\binom{j+1}{k+1}\binom{k}{l}\binom{N-1}{j+2}\binom{-j+N-3}{q}\binom{-l+N-q-3}{s}\binom{-l+N-q-3}{r}\binom{-l+N-q-s-3}{r}r!(-l+N-q-r-s-3)!(s-1)!}{(-l+N-q-2)!(-j+N-1)(N-q-r-s-2)(q+s+1)}$$

$$4S_1(-j+N-1) - 4S_1(-j+N-2) - 2S_1(k)$$

$$-(S_1(-l+N-q-2)+S_1(-l+N-q-r-s-3)-2S_1(r+s))$$

$$+2S_1(s-1)-2S_1(r+s)$$
 + 3 further 6-fold sums

$$\begin{split} \overline{F_0(N)} &= \\ \hline F_0(N) &= \\ \hline \frac{7}{12}S_1(N)^4 + \frac{(17N+5)S_1(N)^3}{3N(N+1)} + (\frac{35N^2-2N-5}{2N^2(N+1)^2} + \frac{13S_2(N)}{2} + \frac{5(-1)^N}{2N^2})S_1(N)^2 \\ &+ (-\frac{4(13N+5)}{N^2(N+1)^2} + (\frac{4(-1)^N(2N+1)}{N(N+1)} - \frac{13}{N})S_2(N) + (\frac{29}{3} - (-1)^N)S_3(N) \\ &+ (2+2(-1)^N)S_{2,1}(N) - 28S_{-2,1}(N) + \frac{20(-1)^N}{N^2(N+1)})S_1(N) + (\frac{3}{4} + (-1)^N)S_2(N)^2 \\ &- 2(-1)^NS_{-2}(N)^2 + S_{-3}(N)(\frac{2(3N-5)}{N(N+1)} + (26+4(-1)^N)S_1(N) + \frac{4(-1)^N}{N+1}) \\ &+ (\frac{(-1)^N(5-3N)}{2N^2(N+1)} - \frac{5}{2N^2})S_2(N) + S_{-2}(N)(10S_1(N)^2 + (\frac{8(-1)^N(2N+1)}{N(N+1)} \\ &+ \frac{4(3N-1)}{N(N+1)}S_1(N) + \frac{8(-1)^N(3N+1)}{N(N+1)^2} + (-22+6(-1)^N)S_2(N) - \frac{16}{N(N+1)}) \\ &+ (\frac{(-1)^N(9N+5)}{N(N+1)} - \frac{29}{3N})S_3(N) + (\frac{19}{2} - 2(-1)^N)S_4(N) + (-6+5(-1)^N)S_{-4}(N) \\ &+ (-\frac{2(-1)^N(9N+5)}{N(N+1)} - \frac{2}{N})S_{2,1}(N) + (20+2(-1)^N)S_{-2,2}(N) + (-17+13(-1)^N)S_{3,1}(N) \\ &- \frac{8(-1)^N(2N+1) + 4(9N+1)}{N(N+1)}S_{-2,1}(N) - (24+4(-1)^N)S_{-3,1}(N) + (3-5(-1)^N)S_{2,1,1}(N) \\ &+ 32S_{-2,1,1}(N) + \left(\frac{3}{2}S_1(N)^2 - \frac{3S_1(N)}{N} + \frac{3}{2}(-1)^NS_{-2}(N)\right) \zeta(2) \end{split}$$

 $F_0(N) =$ $\frac{7}{12}S_{1}(N) = \sum_{i=1}^{N} \frac{1}{i} \frac{1}{N(2N+1)} + (\frac{35N^{2} - 2N - 5}{2N^{2}(N+1)^{2}} + \frac{13S_{2}(N)}{2} + \frac{5(-1)^{N}}{2N^{2}})S_{1}(N)^{2} + (\frac{15N^{2}}{2N^{2}})S_{1}(N)^{2} + (\frac{$ $+ \left(2 + 2(-1)^{N}\right)S_{2,1}(N) - 28S_{-2,1}(N) + \frac{20(-1)^{N}}{N^{2}(N+1)}S_{1}(N) + \left(\frac{3}{4} + (-1)^{N}\right)S_{2}(N)^{2}$ $-2(-1)^{N}S_{-2}(N)^{2} + S_{-3}(N)\left(\frac{2(3N-5)}{N(N+1)} + \left(26 + 4(-1)^{N}\right)S_{1}(N) + \frac{4(-1)^{N}}{N+1}\right)$ $+\left(\frac{(-1)^{N}(5-3N)}{2N^{2}(N+1)}-\frac{5}{2N^{2}}\right)S_{2}(N)+S_{-2}(N)\left(10S_{1}(N)^{2}+\left(\frac{8(-1)^{N}(2N+1)}{N(N+1)}\right)S_{2}(N)+S_{-2}(N)\right)S_{2}(N)+S_{-2}(N)\left(10S_{1}(N)^{2}+\frac{8(-1)^{N}(2N+1)}{N(N+1)}\right)S_{2}(N)$ $+\frac{4(3N-1)}{N(N+1)}S_1(N)+\frac{8(-1)^N(3N+1)}{N(N+1)^2}+\left(-22+6(-1)^N\right)S_2(N)-\frac{16}{N(N+1)}\right)$ $+\left(\frac{(-1)^{N}(9N+5)}{N(N+1)}-\frac{29}{2N}\right)S_{3}(N)+\left(\frac{19}{2}-2(-1)^{N}\right)S_{4}(N)+\left(-6+5(-1)^{N}\right)S_{-4}(N)$ + $\left(-\frac{2(-1)^{N}(9N+5)}{N(N+1)}-\frac{2}{N}\right)S_{2,1}(N) + \left(20+2(-1)^{N}\right)S_{2,-2}(N) + \left(-17+13(-1)^{N}\right)S_{3,1}(N)$ $-\frac{8(-1)^{N}(2N+1)+4(9N+1)}{N(N+1)}S_{-2,1}(N)-(24+4(-1)^{N})S_{-3,1}(N)+(3-5(-1)^{N})S_{2,1,1}(N)$ $+32S_{-2,1,1}(N) + \left(\frac{3}{2}S_{1}(N)^{2} - \frac{3S_{1}(N)}{N} + \frac{3}{2}(-1)^{N}S_{-2}(N)\right)\zeta(2)$

$$\begin{split} \overline{F_0(N)} &= \\ \hline F_0(N) &=$$
$$\begin{split} \overline{F_0(N)} &= \\ \hline F_0(N) &=$$

Example 2:

2-mass 3-loop Feynman integrals



All diagrams are produced with axodraw (J. Vermaseren).



Mellin-Barnesand $_{p}F_{q}$ -technologies expression (95 MB) with

- 150 single sums
- 1000 double sums
- 12160 triple sums
- 1555 quadruple sums



Mellin-Barnesand $_{p}F_{q}$ -technologies expression (95 MB) with

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- 1000 double sums
- 12160 triple sums
- 1555 quadruple sums

Typical triple sum:

$$\begin{split} \sum_{j=0}^{N} \sum_{i=0}^{j} \sum_{k=0}^{i} \frac{(4+\varepsilon)(-2+N)(-1+N)N\pi(-1)^{2-k}}{2+\varepsilon} \times 2^{-2+\varepsilon} e^{-\frac{3\varepsilon\gamma}{2}} \eta^{k} \times \\ & \frac{\Gamma(1-\frac{\varepsilon}{2}-i+j+k)\Gamma(-1-\frac{\varepsilon}{2})\Gamma(2+\frac{\varepsilon}{2})\Gamma(1+N)\Gamma(1+\varepsilon+i-k)\Gamma(-\frac{3\varepsilon}{2}+k)\Gamma(1-\varepsilon+k)\Gamma(3-\varepsilon+k)\Gamma(-\frac{1}{2}-\frac{\varepsilon}{2}+k)}{\Gamma(-\frac{3}{2}-\frac{\varepsilon}{2})\Gamma(\frac{5}{2}+\frac{\varepsilon}{2})\Gamma(2+i)\Gamma(1+k)\Gamma(2-i+j)\Gamma(2-\varepsilon+k)\Gamma(\frac{5}{2}-\varepsilon+k)\Gamma(-\frac{\varepsilon}{2}+k)\Gamma(5+\frac{\varepsilon}{2}+N)} \end{split}$$



Mellin-Barnesand $_{p}F_{q}$ -technologies expression (95 MB) with

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Typical triple sum:

$$\begin{split} \sum_{j=0}^{N} \sum_{i=0}^{j} \sum_{k=0}^{i} \frac{(4+\varepsilon)(-2+N)(-1+N)N\pi(-1)^{2-k}}{2+\varepsilon} \times 2^{-2+\varepsilon} e^{-\frac{3\varepsilon\gamma}{2}} \eta^{k} \times \\ \frac{\Gamma(1-\frac{\varepsilon}{2}-i+j+k)\Gamma(-1-\frac{\varepsilon}{2})\Gamma(2+\frac{\varepsilon}{2})\Gamma(1+N)\Gamma(1+\varepsilon+i-k)\Gamma(-\frac{3\varepsilon}{2}+k)\Gamma(1-\varepsilon+k)\Gamma(3-\varepsilon+k)\Gamma(-\frac{1}{2}-\frac{\varepsilon}{2}+k)}{\Gamma(-\frac{3}{2}-\frac{\varepsilon}{2})\Gamma(\frac{5}{2}+\frac{\varepsilon}{2})\Gamma(2+i)\Gamma(1+k)\Gamma(2-i+j)\Gamma(2-\varepsilon+k)\Gamma(\frac{5}{2}-\varepsilon+k)\Gamma(-\frac{\varepsilon}{2}+k)\Gamma(5+\frac{\varepsilon}{2}+N)} \end{split}$$

6 hours for this sum

 \sim 10 years of calculation time for full expression



 $\underbrace{ \begin{array}{c} \mathsf{Mellin-Barnes-} \\ \mathsf{and} \ _pF_q\text{-technologies} \\ \end{array} }_{p}$

expression (95 MB) with

- 150 single sums
- 1000 double sums
- 12160 triple sums
- 1555 quadruple sums

SumProduction.m (2 hours)

expression (377 MB) consisting of 8 multi-sums



 $\underbrace{ \begin{array}{c} \mathsf{Mellin-Barnes-} \\ \mathsf{and} \ _pF_q\text{-technologies} \\ \end{array} }_{p}$

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expression (377 MB) consisting of 8 multi-sums

EvaluateMultiSums.m

sum	size of sum	summand size of	time of		number of
	(with ε)	constant term	calculation		indef. sums
$\boxed{\sum_{i_4=2}^{N-3} \sum_{i_3=0}^{i_4-2} \sum_{i_2=0}^{i_3} \sum_{i_1=0}^{\infty}}$	17.7 MB	266.3 MB	177529 s	(2.1 days)	1188
$\sum_{i_3=3}^{N-4} \sum_{i_2=0}^{i_3-1} \sum_{i_1=0}^{\infty}$	232 MB	1646.4 MB	980756 s	(11.4 days)	747
$\sum_{i_2=3}^{N-4} \sum_{i_1=0}^{\infty}$	67.7 MB	458 MB	524485 s	(6.1 days)	557
$\sum_{i_1=0}^{\infty}$	38.2 MB	90.5 MB	689100 s	(8.0 days)	44
$\begin{bmatrix} N-3 & i_4-2 & i_3 \\ \sum_{i_4=2}^{N-3} & \sum_{i_3=0}^{N-3} & \sum_{i_2=0}^{N-3} & \sum_{i_1=0}^{N-3} \end{bmatrix}$	1.3 MB	6.5 MB	305718 s	(3.5 days)	1933
$\sum_{i_3=3}^{N-4} \sum_{i_2=0}^{i_3-1} \sum_{i_1=0}^{i_2}$	11.6 MB	32.4 MB	710576 s	(8.2 days)	621
$\sum_{i_2=3}^{N-4} \sum_{i_1=0}^{i_2}$	4.5 MB	5.5 MB	435640 s	(5.0 days)	536
$\sum_{i_1=3}^{N-4}$	0.7 MB	1.3 MB	9017s	(2.5 hours)	68



 $\underbrace{ \begin{array}{c} \mathsf{Mellin-Barnes-} \\ \mathsf{and} \ _pF_q\text{-technologies} \\ \end{array} }_{p}$

expression (95 MB) with

- 150 single sums
- 1000 double sums
- 12160 triple sums
- 1555 quadruple sums

SumProduction.m (2 hours)

expression (377 MB) consisting of 8 multi-sums

> EvaluateMultiSums.m (3 month)

expression (154 MB) consisting of 4110 indefinite sums

Most complicated objects: generalized binomial sums, like

$$\sum_{h=1}^{N} 2^{-2h} (1-\eta)^h \binom{2h}{h} \left(\sum_{i=1}^{h} \frac{2^{2i} (1-\eta)^{-i}}{i\binom{2i}{i}} \right) \left(\sum_{i=1}^{h} \frac{(1-\eta)^i \binom{2i}{i}}{2^{2i}} \right) \times \left(\sum_{i=1}^{h} \frac{2^{2i} (1-\eta)^{-i} \sum_{j=1}^{i} \frac{\sum_{i=1}^{j} \frac{(1-\eta)^k}{k}}{j}}{i\binom{2i}{i}} \right) \right)$$



 $\underbrace{ \mathsf{Mellin-Barnes-}}_{\text{and } pF_q-\text{technologies}} \xrightarrow{}$

expression (95 MB) with

- 150 single sums
- 1000 double sums
- 12160 triple sums
- 1555 quadruple sums

SumProduction.m (2 hours)

expression (377 MB) consisting of 8 multi-sums

> EvaluateMultiSums.m (3 month)

expression (8.3 MB) consisting of 74 indefinite sums

Sigma.m (32 days)

expression (154 MB) consisting of 4110 indefinite sums

Evaluation of Feynman Integrals



Evaluation of Feynman Integrals

